

*Semi-formal models to support program development:
autonomic management within component based
parallel and distributed programming*

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FMCO 2008

Sophia Antipolis - October 2008

Contents

- Introduction
 - Functional vs. non-functional concerns
 - Autonomic management of non-functional concerns
- “Semi-formal” handling of non-functional concerns
 - ORC
- Use case
 - Performance tuning in stream parallel component compositions
- Conclusions

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Functional & non functional concerns

- Functional
 - all aspects related to **what** is computed
- Non-functional
 - all aspects related to **how** the result is computed
- In parallel distributed programming
 - *functional*: the algorithm, the kind of parallel pattern used
 - *non-functional*: parallelism degree, load balance, fault tolerance, security, ...

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Application
programmer concerns

Functional & non functional concerns

□ Functional

- all aspects related to **what** is computed

Application

programmer concerns

□ Non-functional

- all aspects related to **how** the result is computed

System programmer concerns

□ In parallel distributed programming

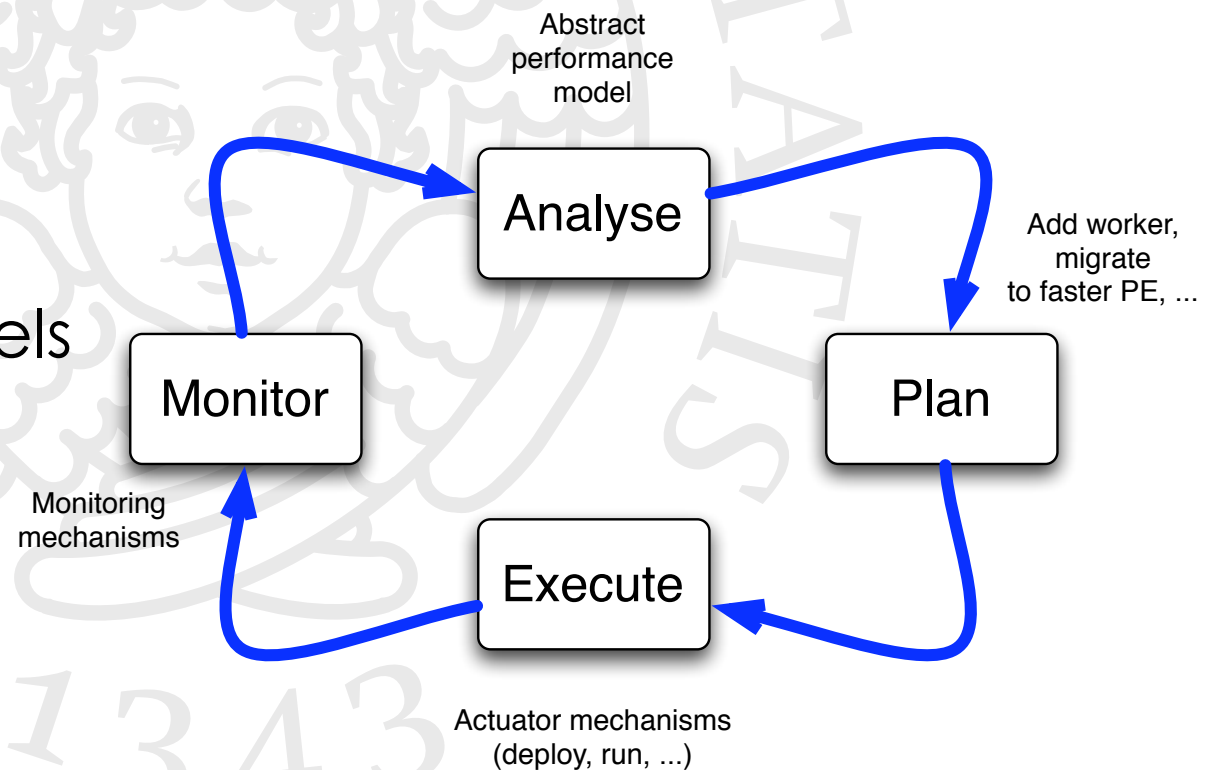
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Autonomic management of non functional concerns

- Autonomic management
 - control loop: monitor → analyze → plan → execute
 - monitoring
 - ◆ mechanisms
 - analyzing
 - ◆ reference models
 - planning
 - ◆ strategies
 - executing
 - ◆ mechanisms

Autonomic management of non functional concerns

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Reactive autonomic management

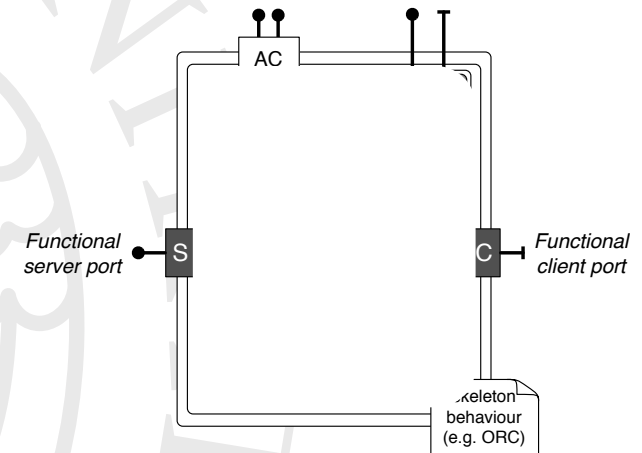
- Monitoring
 - non-invasive, immediate, effective
- Analysis
 - automatic, prioritized, extendible
- Planning
 - target architecture specific, optimized
- Execute
 - efficient, fast

Autonomic management in GCM

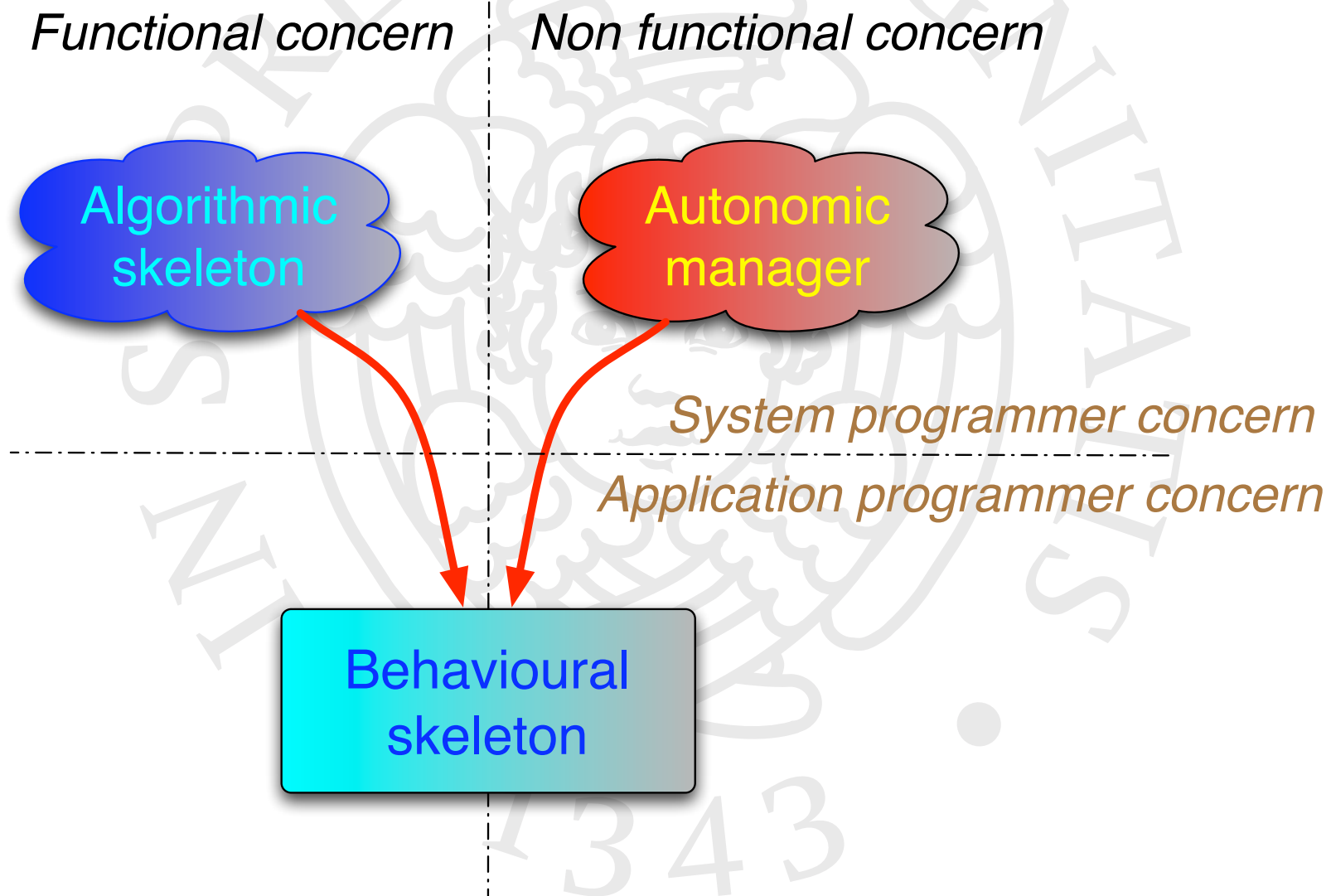
- Behavioural skeleton concept
 - co-design of
 - ◆ parallelism exploitation pattern
 - ◆ autonomic management
- Performance management
 - initial setup (parallelism degree)
 - optimization/tuning (load balancing, fault tolerance)
 - user driven (contract/SLA)

Autonomic management in GCM

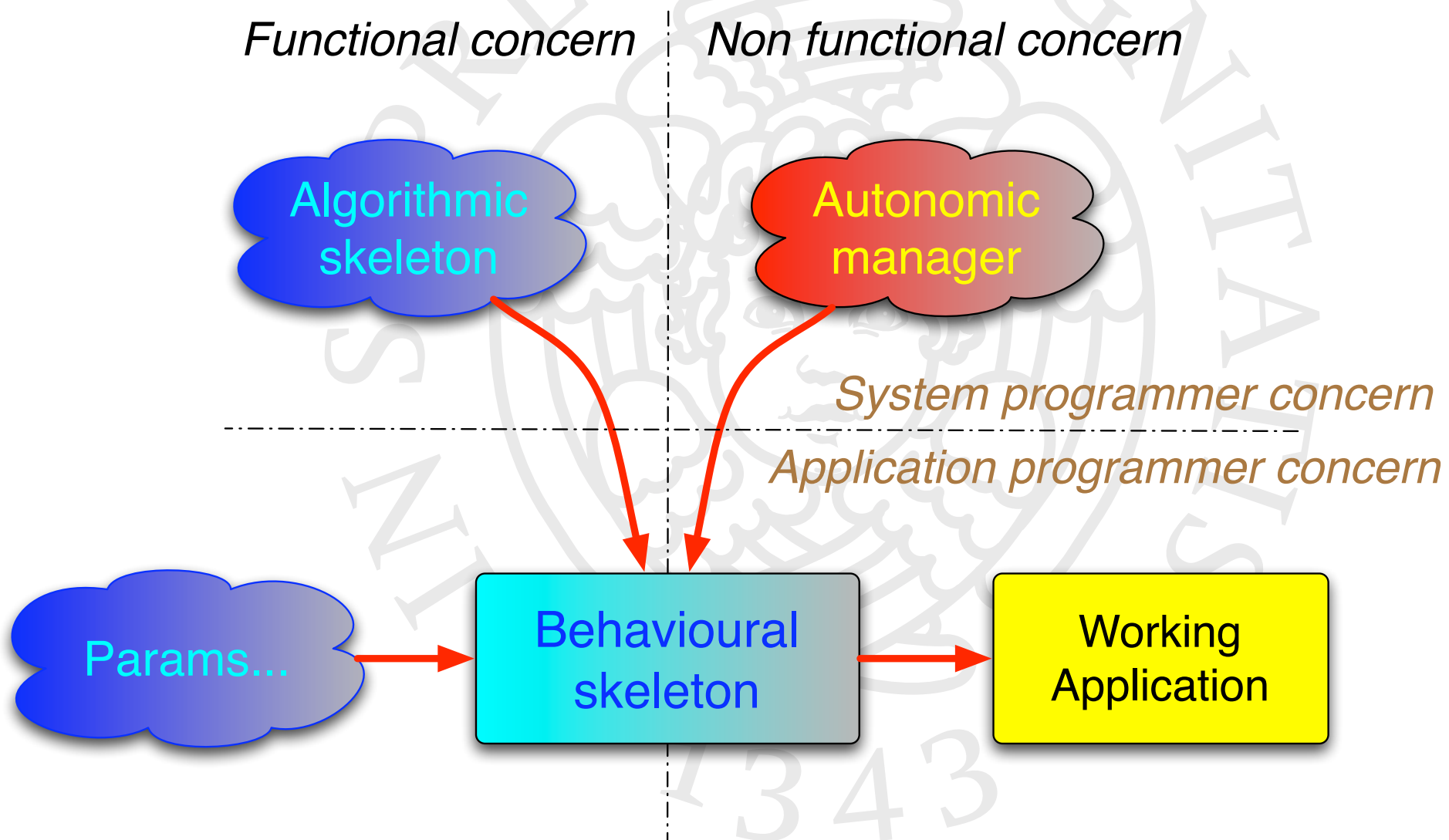
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Behavioural skeleton

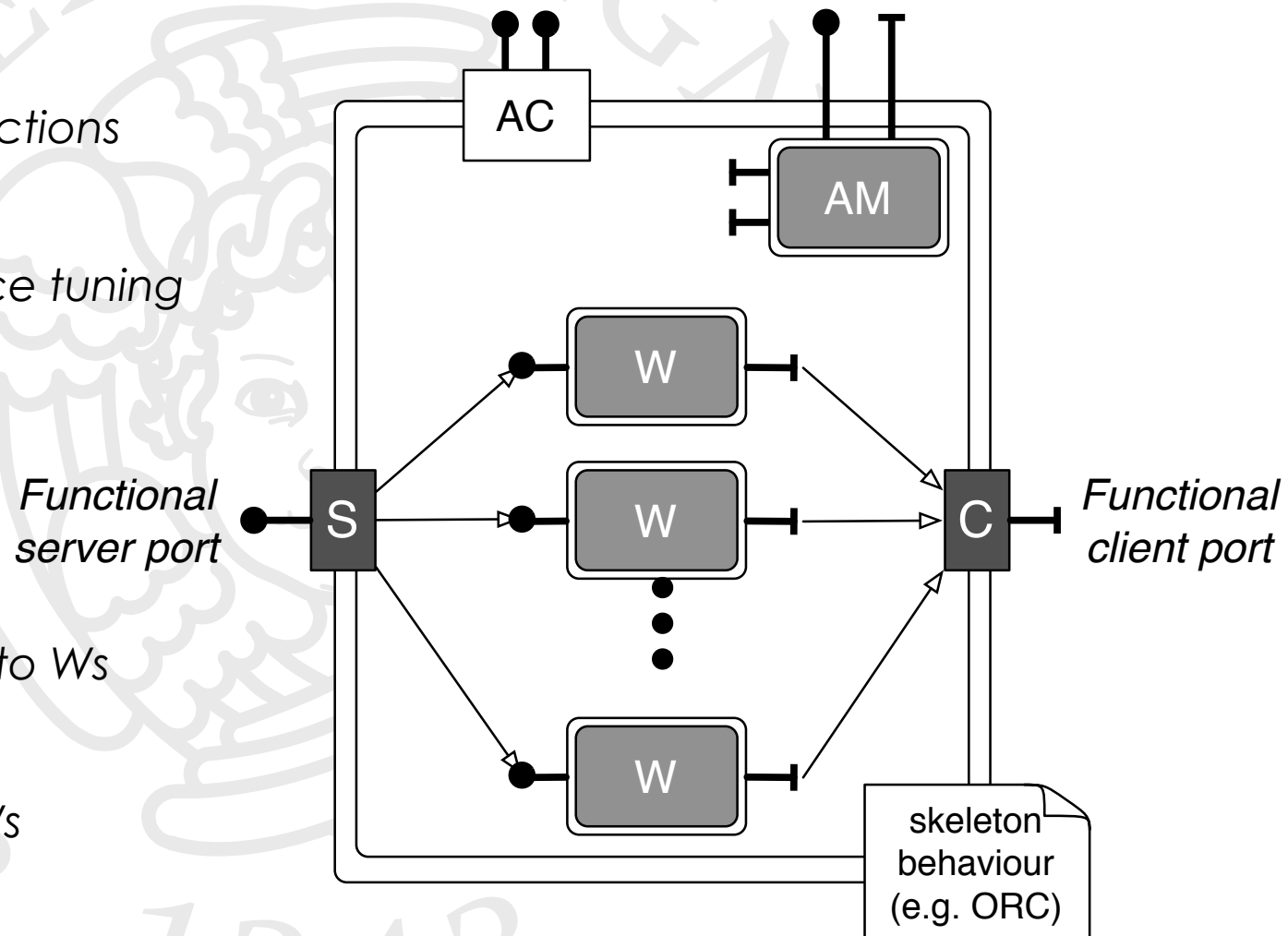


Behavioural skeleton



Behavioural skeleton sample: functional replication

- Autonomic Controller
implements passive actions
- Autonomic Manager
manages performance tuning
- S port
distributes input tasks to Ws
- C Port
collects results from Ws
- Worker components
compute results (functional)



Rule based autonomic management

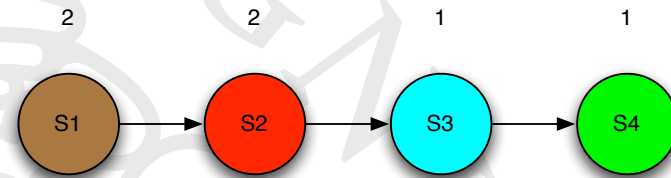
- triggering of actions
 - first order logic formulae over monitoring figures
- analysis
 - ordering / scheduling of the *fired* triggers
- planning/execute
 - sequence of mechanism invocation
- GCM uses JBoss rules (drools)
 - *rulename* salience *nn* when ... then ...

Already implemented (single pattern/ manager)

- parallelism degree adjustment
 - increase
 - ◆ suitable input pressure & unsatisfied contract
 - decrease
 - ◆ over satisfied contract
 - ◆ unsuitable input pressure
- fault tolerance
 - ◆ automatic recovery of faulty resources (muskel)

To be implemented (hierarchical pattern/ manager)

- change in (nesting of) parallelism exploitation patterns used



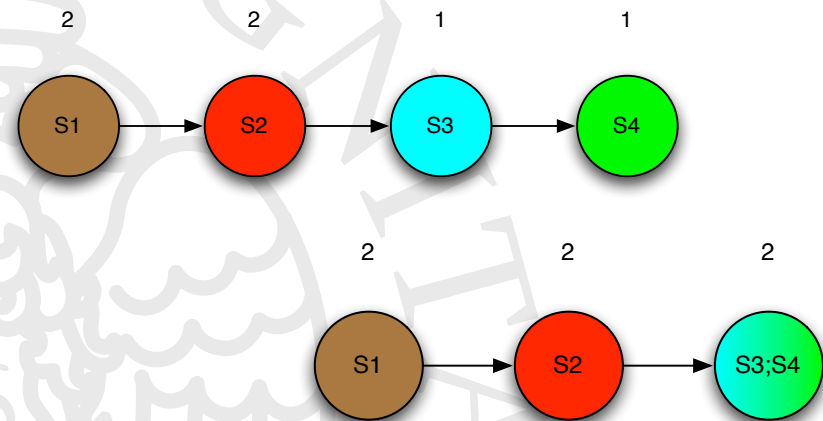
- pipeline stage unbalance
 - ◆ stage collapsing
 - ◆ farm out stage
 - ◆ combination of collapsing and farming out

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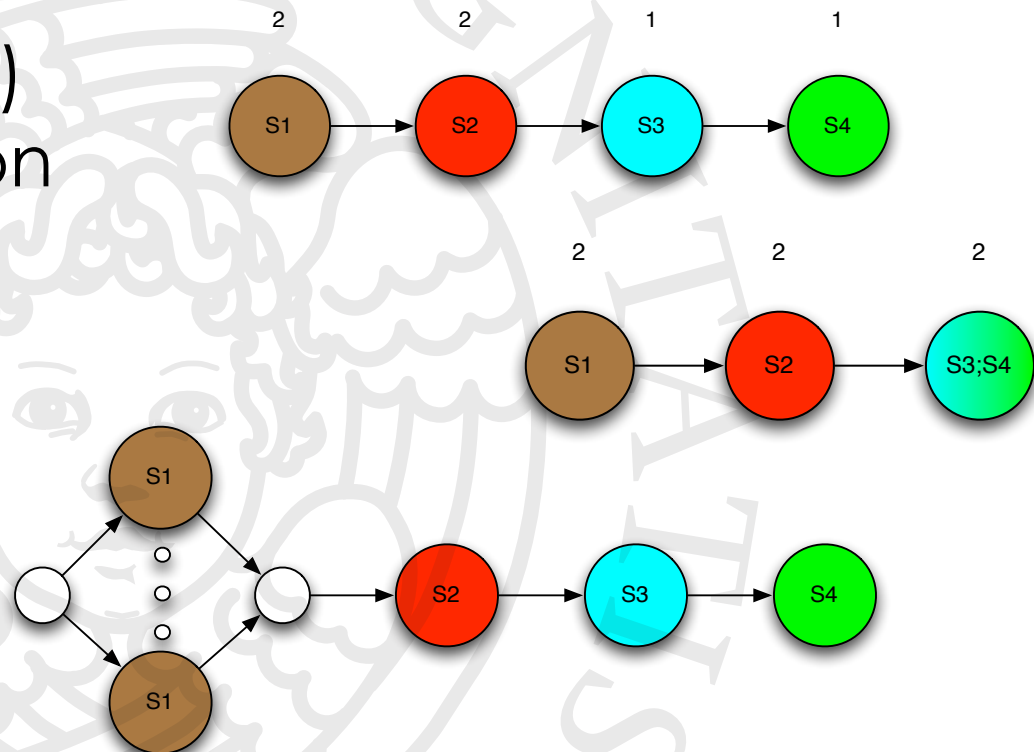


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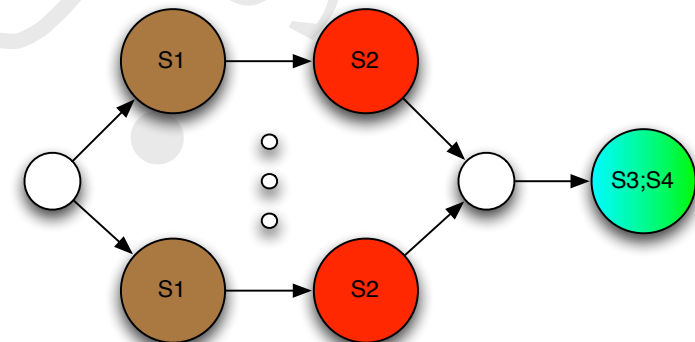
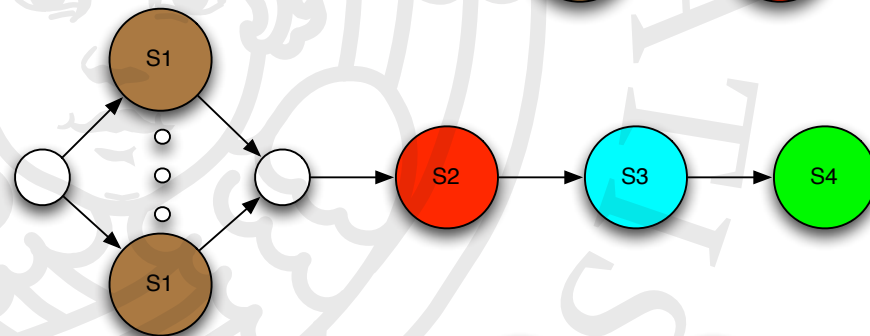
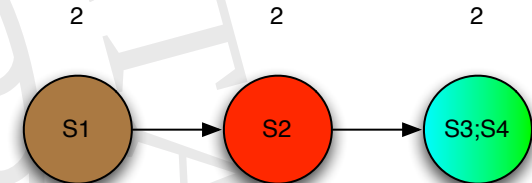
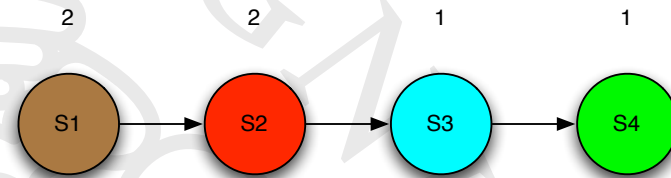
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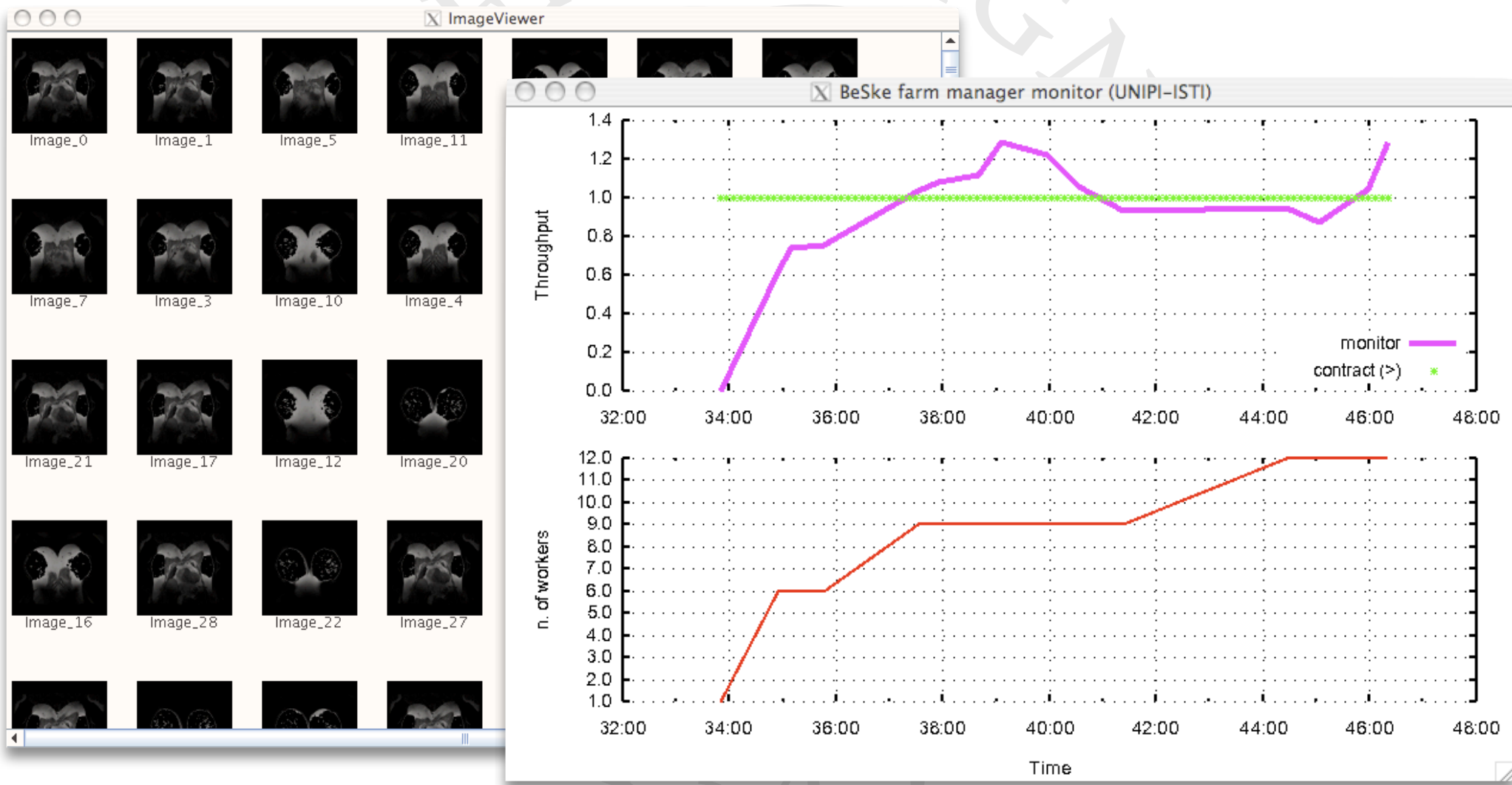
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Autonomic performance management @ work



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Tools to support reasoning about autonomic management

- the two extremes
 - formal tools
 - ◆ consistent background needed
 - ◆ nice results demonstrated
 - possibly with limited scope
 - implementation
 - ◆ consistent background & ability needed
 - ◆ nice results “demonstrated”
 - possibly requiring a huge amount of time

Semi-formal tools

- preserve part of the knowledge typical of programmers
- preserve part of the techniques typical of formal tools
- e.g.
 - a framework
 - ◆ suitable to model interesting properties of a (distributed/parallel) program
 - ◆ synthetic
 - ◆ supporting program transformations

Orc Language Project

[Home](#)[Download](#)[Documentation](#)[Research](#)[Community](#)

Orc is ...

A novel language for distributed and concurrent programming. Orc provides uniform access to computational services, including distributed communication and data manipulation, through *sites*. Using three simple concurrency primitives, the programmer *orchestrates* the invocation of sites to achieve a goal, while managing timeouts, priorities, and failures.

- ❑ Introduced by Misra and Cook in early '00
- ❑ provides primitive *combinators* for parallelism and non determinism (asymmetric parallelism)
- ❑ look&feel close to a programming language

What is it for?

- Concurrent and distributed programming
- Workflows (business process automation)
- Discrete event simulation
- Web service mashups

ORC in a slide

- site : local or remote (unreliable) unit of computation
- combinators
 - $a \mid b$ site or expression a and b evaluated in parallel
 - $a \gg b$ (or $a \succ x \succ b$) a evaluated first, then b
 - $f(x)$ where $x \in (a \mid b)$ a, b and f started, as soon as either a or b produce a value, it is bound to x and $(a \mid b)$ is terminated
(new syntax $f(x) \langle x \langle (a \mid b) \rangle$)
- Functions
 - $\text{def } f(\text{param}) = \dots$
- predefined sites + channels
 - $\text{if}, \text{Rtimer}, +, -, \dots, \text{ch.get}(), \text{ch.put}(x)$

most/all the
abstractions needed
are there

Sample usage of ORC

- Reverse engineering (modelling) of muskel, a full Java/RMI skeleton library maintained at University of Pisa

```

system(pgm, tasks, contract, G, t)  $\triangleq$ 
  taskpool.add(tasks) | discovery(G, pgm, t) | manager(pgm, contract, t)
discovery(G, pgm, t)  $\triangleq$  (|g ∈ G ( if remw ≠ false  $\gg$  rworkerpool.add(remw)
  where remw :∈
    ( g.can_execute(pgm) | Rtimer(t)  $\gg$  let(false) )
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  |i : 1 ≤ i ≤ contract : (rworkerpool.get > remw > ctrlthreadi(pgm, remw, t))
  | monitor
ctrlthreadi(pgm, remw, t)  $\triangleq$  taskpool.get > tk >
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triggering

monitoring

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monitor  $\triangleq$  alarm.get > i > rworkerpool.get > remw > ci.put(remw)
   $\gg$  monitor
  
```

plan
execute

triggering

monitoring

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M. Aldinucci, M. Danelutto, and P. Kilpatrick. Management in distributed systems: a semi-formal approach. In A. M. Kermarrec, T. Priol, and L. Bougé, editors, Proc. of 19th Intl. Euro-Par 2007 Parallel Processing, volume 4641 of LNCS, pages 651–661, Rennes, France, August 2007. Springer

```

    discovery(G, pgm, t) ≜ (|g ∈ G ( if remw ≠ false ≫ rworkerpool.add(remw)
    where remw :∈
    ( g.can_execute(pgm) | Rtimer(t) ≫ let(false) )
    )
    ) ≫ discovery(G, pgm, t)

    manager(pgm, contract, t) ≜
    |i : 1 ≤ i ≤ contract : (rworkerpool.get > remw > ctrlthread_i(pgm, remw, t))
    | monitor

    ctrlthread_i(pgm, remw, t) ≜ taskpool.get > tk >
    ( -if valid ≫ resultpool.add(r) ≫ ctrlthread_i(pgm, remw, t)
    | if ¬valid ≫ ( taskpool.add(tk)
    | alarm.put(i) ≫ c_i.get > w ≫ ctrlthread_i(pgm, w, t)
    )
    )

    where (valid, r) :∈
    ( remw(pgm, tk) > r > let(true, r) | Rtimer(t) ≫ let(false, 0) )

    monitor ≜ alarm.get > i > rworkerpool.get > remw > c_i.put(remw)
    ≫ monitor
    
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ORC for autonomic management

- semi formal derivation / proof of rewriting rules
- e.g. $\text{pipe}(f,g) = \text{pipe}(\text{farm}(f),g)$
 - simple example, technique to be used for more complex derivations
 - first step: formalization of skeletons
 - second step: semi formal processing of ORC expressions

Formalization of skeletons

$$\text{pipe}(A, B, ch_{in}, ch_{out}) = \text{stage}(A, ch_{in}, ch_{new}) \\ | \text{stage}(B, ch_{new}, ch_{out})$$
$$\text{stage}(A, ch_{in}, ch_{out}) = ch_{in}.get() > task > A(task) > \\ result > ch_{out}.put(result) >> \\ \text{stage}(A, ch_{in}, ch_{out})$$
$$\text{farm}(W, nw, c_{in}, ch_{out}) = | i = 1, nw : \text{Worker}_i(W, c_{in}, ch_{out})$$
$$\text{Worker}(W, c_{in}, ch_{out}) = ch_{in}.get() > task > W(task) > \\ result > ch_{out}.put(result) >> \\ \text{Worker}(W, c_{in}, ch_{out})$$

Semi-formal processing of ORC expressions

- e.g. exploit semantics for channels
 - matching get/put pair collapsing
 - in actual traces
 - (R free for x !)

$$\frac{(a > x > ch.put(x) > R) \mid (\dots >> ch.get() > y > S)}{R \mid \dots >> a > y > S}$$

Sample semi-formal reasoning

$\text{pipe}(A, B, c_1, c_3) = \text{stage}(A, c_1, c_2) \mid \text{stage}(B, c_2, c_3)$

$\text{seq}(A, B, ch_{in}, ch_{out}) = c_{in}.get() > x > A(x) > y > B(y) > z > ch_{out}.put(z)$

Sample semi-formal reasoning

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$c_1.get() > t > A(t) > y > c_2.put(y) \gg \text{stage}(A, c_1, c_2)$

Sample semi-formal reasoning

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$c_1.get() > t > A(t) > y > c_2.put(y) >> \text{stage}(A, c_1, c_2)$

$c_2.get() > t > B(t) > y > c_3.put(y) >> \text{stage}(B, c_2, c_3)$

Sample semi-formal reasoning

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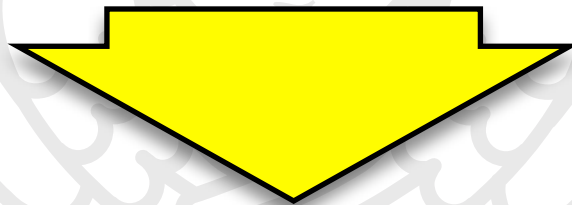
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Sample semi-formal reasoning

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$$c_2.get() > t > B(t) > y > c_3.put(y) >> \text{stage}(B, c_2, c_3)$$


$$c_1.get() > x > \text{stage}(\text{seq}(A, B))(x) > y > c_3.put(y)$$

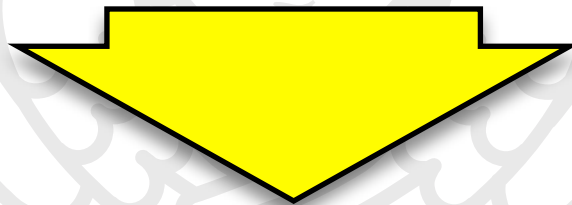
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$$c_2.get() > t > B(t) > y > c_3.put(y) >> \text{stage}(B, c_2, c_3)$$



$$c_1.get() > x > \text{stage}(\text{seq}(A, B))(x) > y > c_3.put(y)$$

$$\text{pipe}(A, B, ch_{in}, ch_{out}) \leftrightarrow \text{seq}(A, B, ch_{in}, ch_{out})$$

Design of autonomic managers

$$BSekl(Sk, Mgr, SLA) = Sk \mid Mgr(Sk, SLA)$$

$$\begin{aligned}
 Mgr(Sk, SLA) = & \text{distribute}(Sk, SLA) > s > \\
 & \text{monitor}(s) > m > \\
 & \text{analyse}(s, m) > (b, p, v) > \\
 & ((if(b) >> \text{adapt}(s, p) > s1 > \\
 & \quad Mgr(s1, SLA)) \\
 & \mid (if(\sim b) >> \text{raise}(v) > \\
 & \quad Mgr(s, \text{passiveMode}(SLA))))
 \end{aligned}$$

Design of autonomic managers

$$BSekl(Sk, Mgr, SLA) = Sk \mid Mgr(Sk, SLA)$$

$$Mgr(Sk, SLA) = \text{distributed} \left(\begin{array}{l} \text{mode}(s) > SLA > s > \\ \text{mode}(s) > m > \\ \text{raise}(s, m) > (b, p, v) > \\ ((if(b) >> adapt(s, p) > s1 > \\ Mgr(s1, SLA)) \\ | (if(\sim b) >> raise(v) > \\ Mgr(s, passiveMode(SLA)))) \end{array} \right)$$

Adaptation in BS

$$\begin{aligned} \text{adapt}(\text{pipe}(A, \text{pipe}(B, C)), \text{plan}) = & \\ & (\text{if}(\text{plan} = \text{collapseFirst}) \gg \text{pipe}(\text{seq}(A, B), C)) \\ & | (\text{if}(\text{plan} == \text{collapseLast}) \gg \text{pipe}(A, \text{seq}(B, C))) \\ & | (\text{if}(\text{plan} == \text{farmoutFirst}) \gg \text{pipe}(\text{farm}(A), \text{pipe}(B, C))) \end{aligned}$$

- modelling of management *before* actual implementation

Adding metadata

- annotations on Orc code
 - modelling several non functional concerns
 - e.g. security, communication costs

- formal process deriving the aggregated metadata from primitive/elementary/ground one (synthesis) or primitive metadata from aggregated (analysis)

Sample metadata: process placement

- Analysis
 - placement annotations
 - policy managing nested skeleton annotations

$$\frac{\textit{placement}(\textit{pipe}(A, B), \textit{loc}(X)) \wedge \textit{distribPolicy}(\textit{keep})}{\textit{placement}(A, \textit{loc}(X)) \wedge \textit{placement}(B, \textit{loc}(X))}$$

$$\frac{\textit{placement}(\textit{pipe}(A, B), \textit{loc}(X)) \wedge \textit{distribPolicy}(\textit{distrib})}{\textit{placement}(A, \textit{loc}(\textit{fresh}())) \wedge \textit{placement}(B, \textit{loc}(\textit{fresh}()))}$$

Sample metadata usage

- skeleton program + placement metadata (includes support: channels, manager process(es), ...)
- communication cost

$$\frac{\textit{placement}(ch.get(), loc(X) \wedge \textit{placement}(ch, loc(Y)))}{\textit{nonLocalCost}(ch.get())}$$

$$\frac{\textit{placement}(ch.get(), loc(X) \wedge \textit{placement}(ch, loc(Y)))}{\textit{nonLocalCost}(ch.get())}$$

- automatic derivation of communication cost in typical traces of execution

Sample metadata usage

- skeleton program + placement metadata / support: channels, manager processes
- communication cost

$placement(ch, loc(X))$

$cost(ch.get())$

$ch.get(), loc(X) \wedge placement(ch, loc(Y))$

$nonLocalCost(ch.get())$

- automatic derivation of communication cost in typical traces of execution

Sample metadata: security

- Synthesis

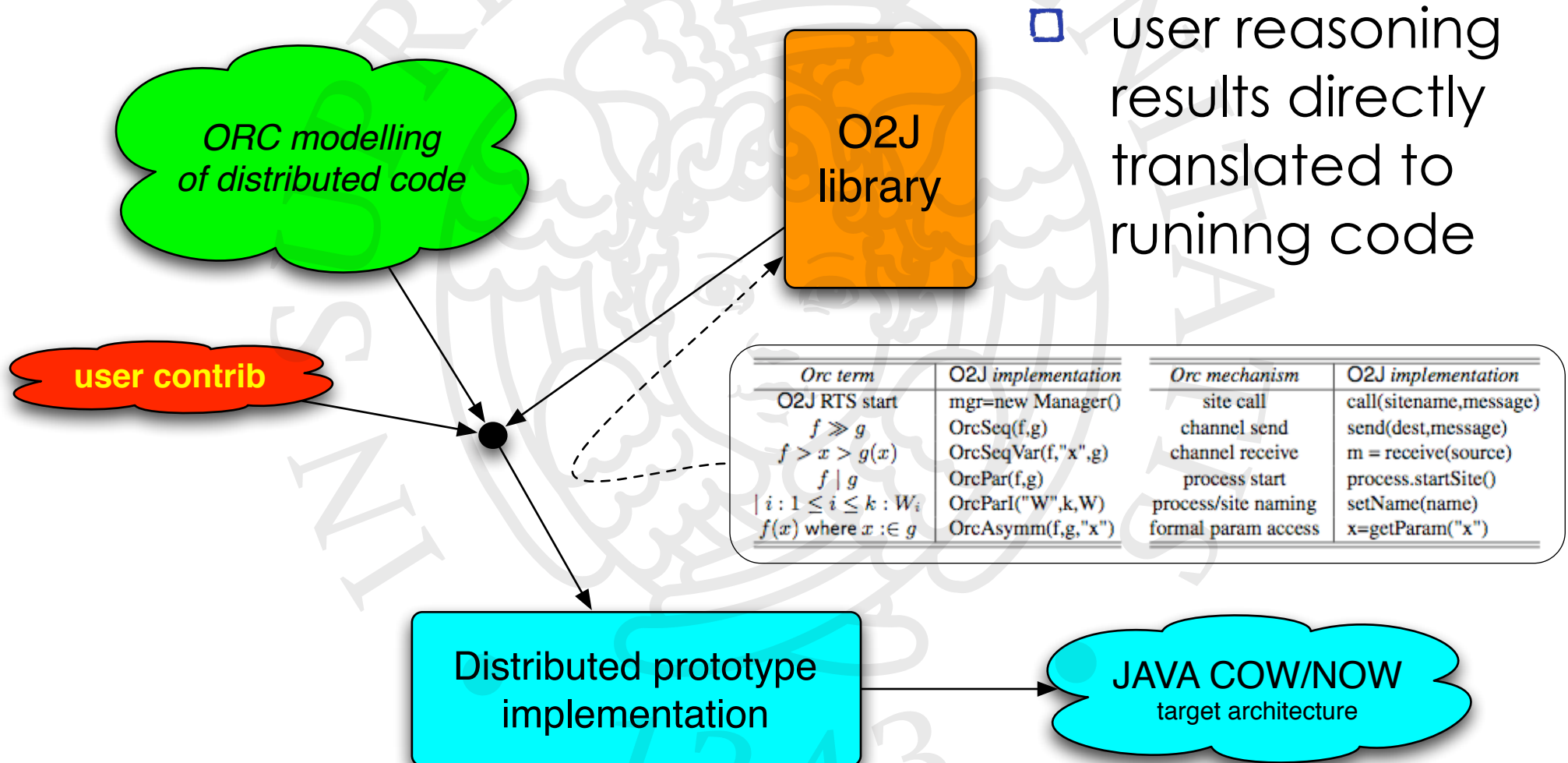
$$\frac{\textit{placement}(A, \textit{loc}(X)) \wedge \textit{insecure}(X)}{\textit{insecure}(\textit{pipe}(A, _)) \textit{insecure}(\textit{pipe}(_, A))}$$

- marking of root depending on the marks at leaves

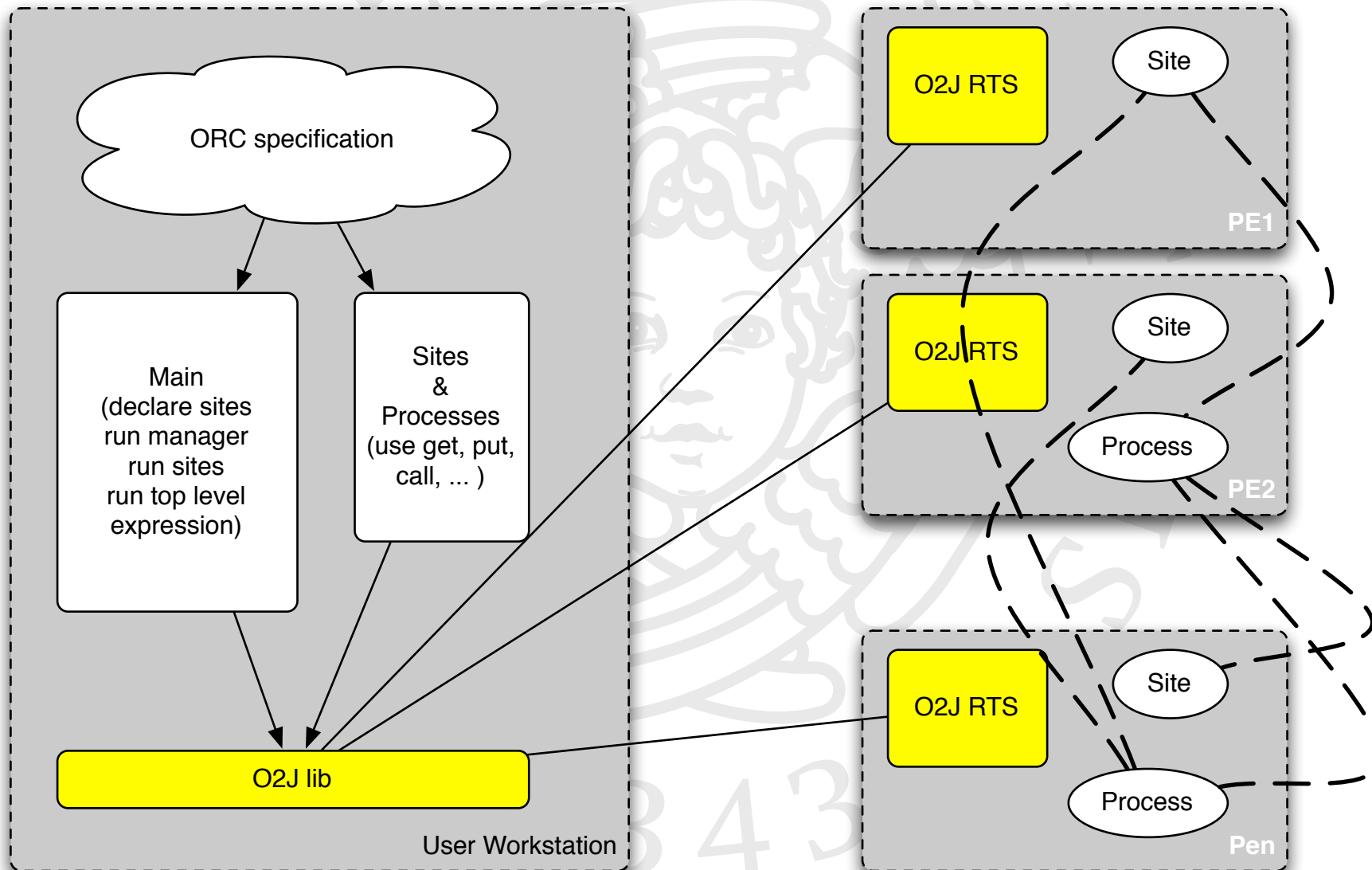
- Usage:

- node marking => securing code and data only when needed

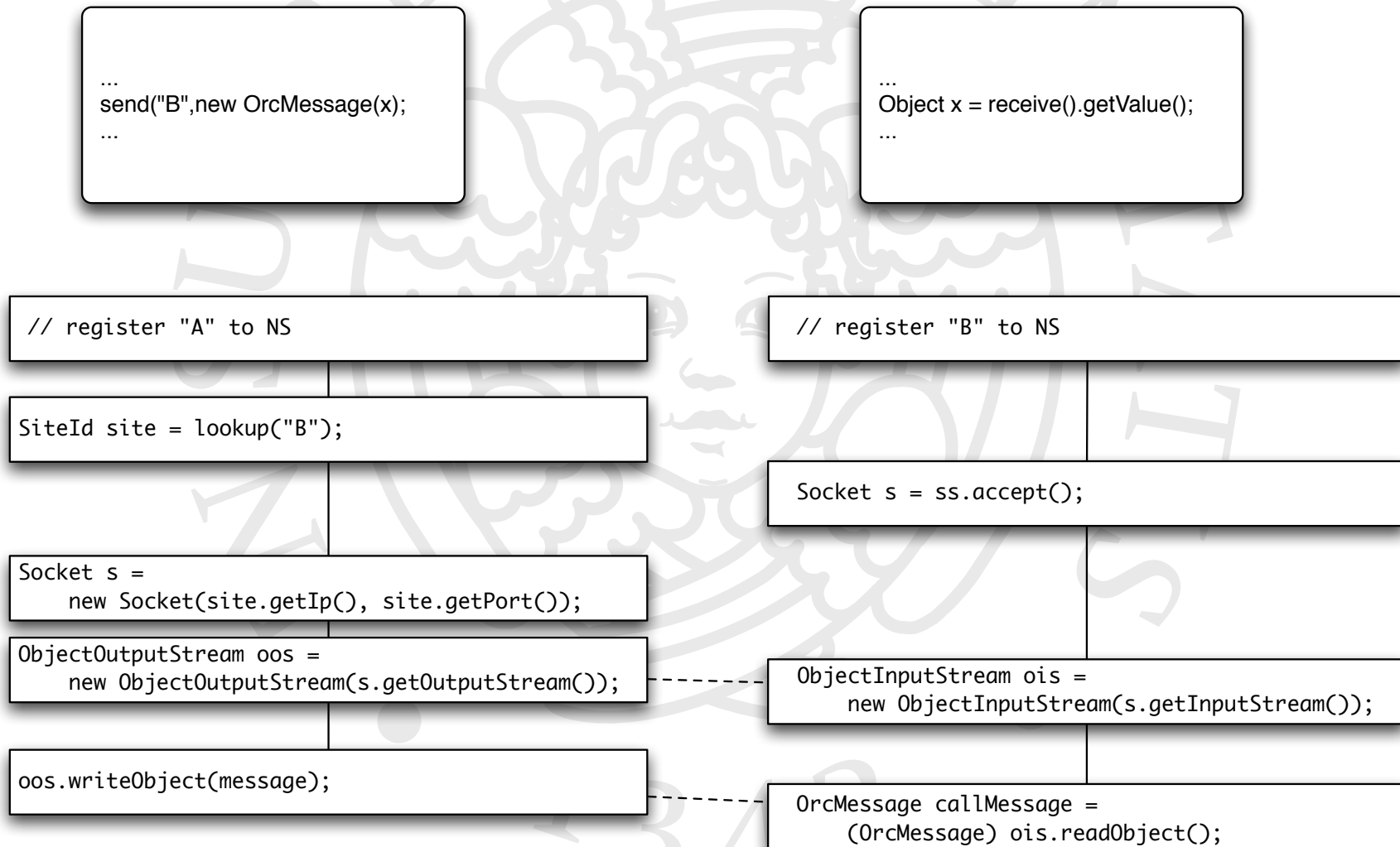
An Orc based development framework



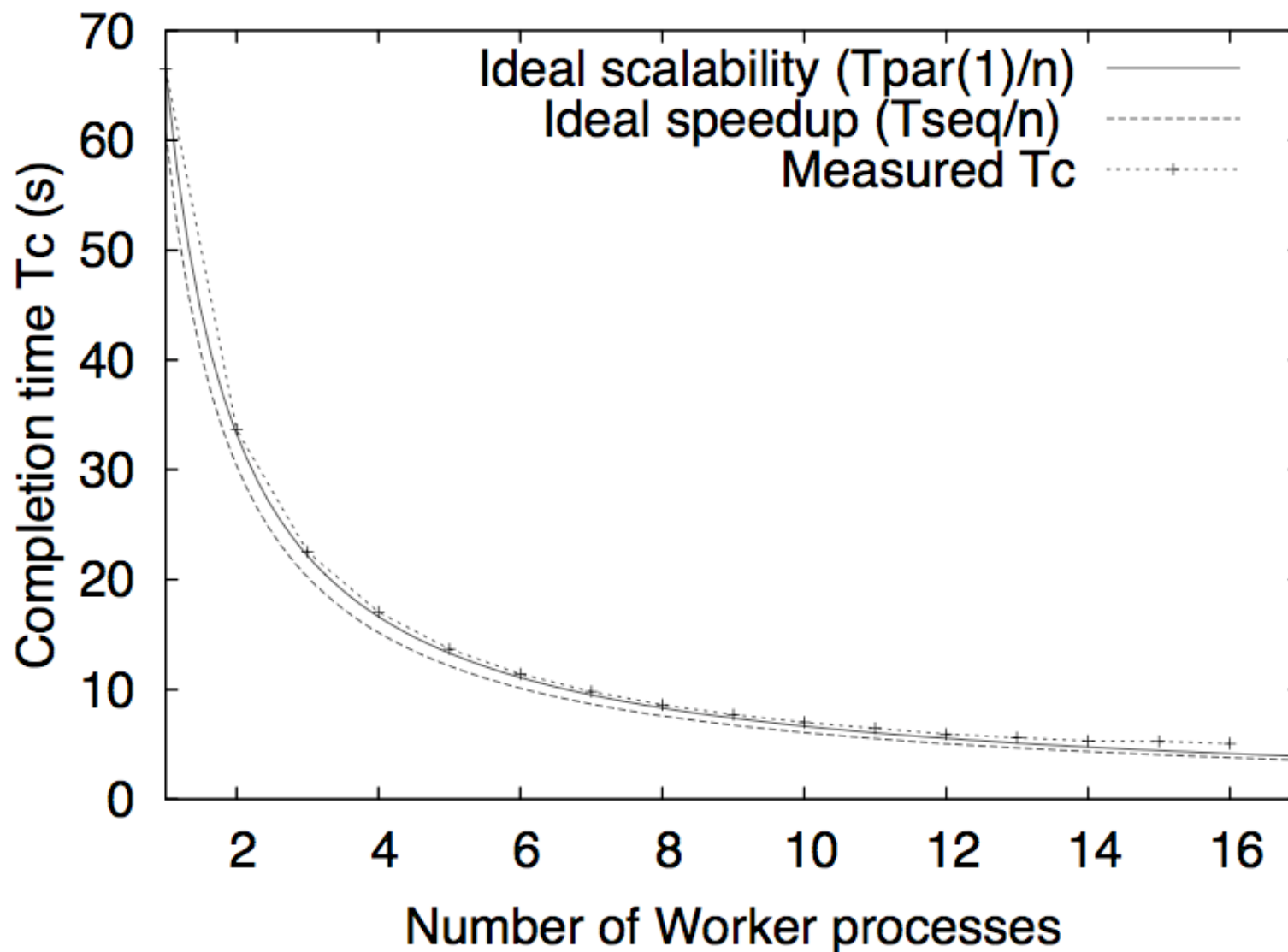
O2J internals (abstract view)



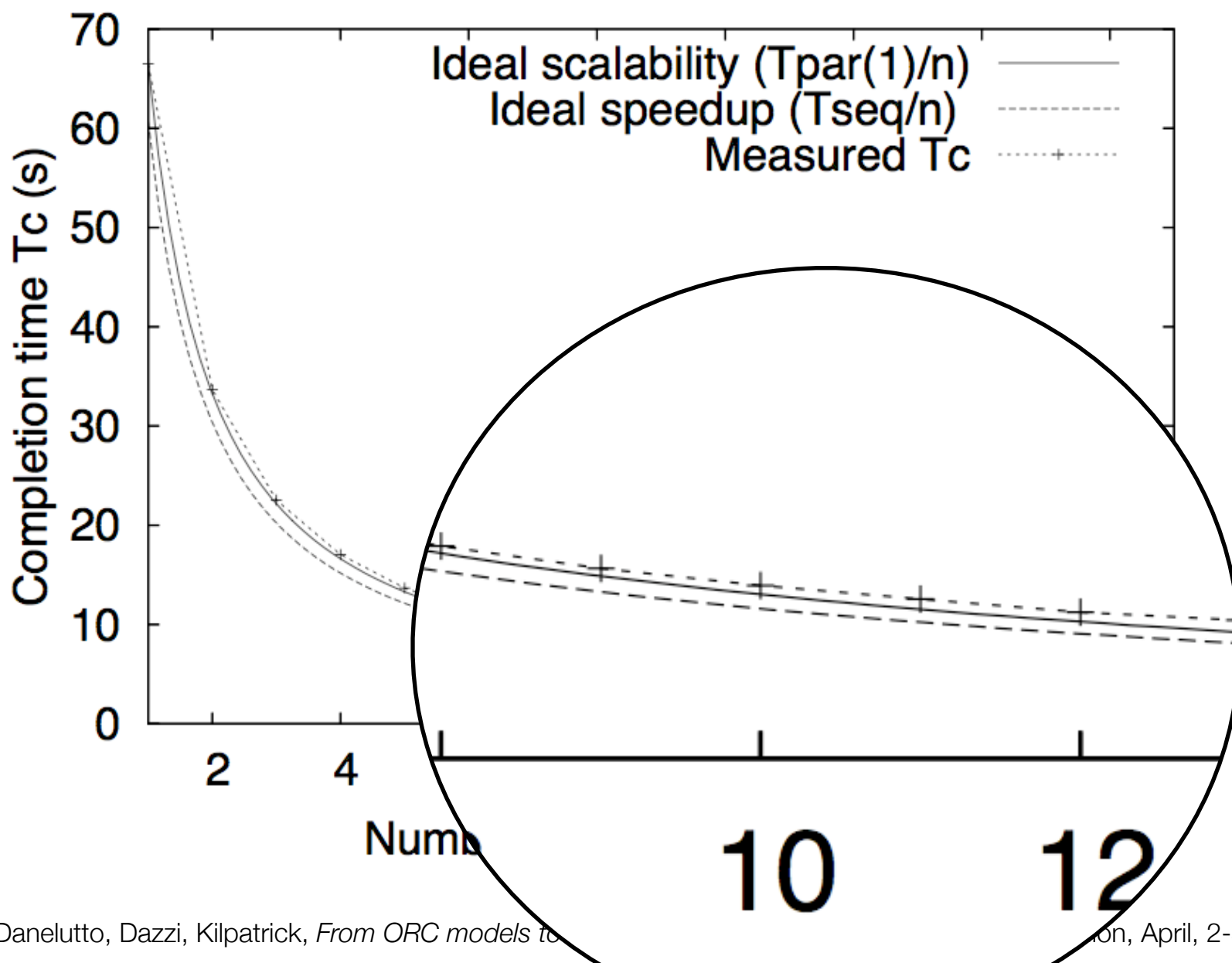
O2J internals (2)



Experimental results



Experimental results



Conclusions

- Autonomic management of non functional features
 - a must
 - a complex task
- Semi formal modelling
 - provides insights and design hints
 - can be used to support reasoning
 - ◆ event stronger with proper metadata
- Experience in GCM / CoreGRID / GridCOMP

Thank you for your attention

Any questions ?