





CoreGRID: European Research Network on Foundations, Software Infrastructures and Applications for large scale distributed, GRID and Peer-to-Peer Technologies

Autonomic Components in GCM

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Outlime

Activities held in - CoreGRID Institute on Programming Models - GridCOMP spin-off project (STREP)

Part I (assessed work)

- Motivations
 - # GCM (coreGrid Component Model)
 - * why adaptive and autonomic management, why skeletons
- behavioural skeletons (in insulation)
 - # demo

Part II (ongoing and future work)

- # formalisation of component and services
- component and service is not a dichotomy
- ** static and dynamic adaptation **should** not be a dichotomy





CGM MODEL KEY POINTS

#Hierarchic model

- expressiveness
- structured composition

Interactions among components

- collective/group
- configurable/programmable
- not only RPC/RMI, but also stream/event

Non-Functional aspects and QoS control

- autonomic computing paradigm
 - adaptive and autonomic components





WHY AUTONOMIC COMPUTING

%// programming & the grid

- concurrency exploitation, concurrent activities set up, mapping/scheduling, communication/synchronisation handling and data allocation, ...
- manage resources heterogeneity and unreliability, networks latency and bandwidth unsteadiness, resources topology and availability changes, firewalls, private networks, reservation and jobs schedulers, ...

... and a non trivial user-defined QoS for applications not easy leveraging only on middleware

our approach: high-level methodologies + tools





WHY AUTONOMIC COMPUTING (USER-DEFINED QOS REQUIREMENTS FOR APPS) Performance

- the app should sustain x transactions per second
- the app should complete each transaction in t seconds

Security

- the link between P1 and P2 should be secured with k-strong encryption
- $^{\circ}$ the DB service is exposed by platform P3

Fault-tolerance

 \circledast the parallel server should survive to the failure of *y* platforms

... then consider that x, t, P1, P2, P3, k, y can dynamically change as may dynamically change the performance and the state of the running environment





WHY SKELETONS

Management is difficult

- application change along time (ADL not enough)
 - how "describe" functional, non-functional features?
- the low-level programming of component and its management is simply too complex

Component reuse is already a problem

- specialising component yet more with management strategy would just worsen the problem
- especially if the component should be reverse engineered to be used (its behaviour may change along the run)





BEHAVIOURAL SKELETONS IDEA

- Represent an evolution of the algorithmic skeleton concept for component management
 - abstract parametric paradigms of component assembly
 - specialised to solve one or more management goals
 - self-configuration/optimization/healing/protection.
 - carry a semi-formal/formal description and an implementation
 - they are component factories, actually
- Are higher-order components
- Are not exclusive
 - an be composed with non-skeletal assemblies via standard components connectors
 - overcome a classic limitation of skeletal systems



ABC = Autonomic Behaviour Controller (implements mechanisms) AM = Autonomic Manager (implements policies) B/LC = Binding + Lifecycle Controller

Part I: BeSke (in insulation) FUNCTIONAL REPLICATION (GCM IMPLEMENTATION)

1. Choose a schema e.g. functional replication ABC API is chosen accordingly

- 2. Choose an inner component compliant to BeSke constraints
- 3. Choose behaviour of ports e.g. unicast/from_any, scatter/gather
- 4. **Run** your application *then trigger adaptations*
- 5. Automatise the process with a Manager











EXAMPLE: DATA PARALLEL (STATELESS)





DATA PARALLEL (STATEFUL, DISTR. STATE)



nformation Society

Notes

- any number of server and client ports (either RPC or stream, in theory)
- the model cannot (structurally) enforce init happens before requests on other ports
- port reconfiguration and data redistribution should be atomic (no tasks should be distributed in the middle.
- data can be reconfigured in a distributed way (provided a suitable data port abstraction is defined)





VARIATIONS AND FLAVOURS (EXAMPLES)



Functional Replication

- Farm/parameter sweep (self-optimization)
- Data-Parallel (self-configuring map-reduce)
- Active/Passive Replication (self-healing)

Proxy

Pipeline (coupled self-protecting proxies)

Wrappers

Facade (self-protection)

















P3 P3•P4 P4

PART II

manager



P2

* formalisation of component and services

- adaptations, QoS contracts, orchestration of managers (as services)
- component and service is not a dichotomy
 - # from GCM/Proactive to SCA/Tuscany
- static and dynamic adaptation should not be a dichotomy
 - * if we care about performance

P5

P6



MANAGER FORMALISATION & DESIGN

- # Hierarchic assemblies of component that may structurally change at run-time. Issues:
 - Formally represent adaptations
 - they should be described in the AM and automatically applied
 - the ADL give just a static view of the assembly
 - Formally represent QoS contracts
 - they should be described in the AM and automatically evaluated
 - they should be projected and joint (almost automatically)
 - Describe the interaction/orchestration among managers
 - Globally, managers describe a distributed algorithm
- Some hints presented here
 - … but still many open problems (just a few discussed here)





FORMALISING ADAPTATIONS

Graphs + graph rewriting

- rewriting rules represent possible adaptation patterns
 - enough expressive ... even too much
 - some formalisation do not capture important concepts for // computing such as locality of the rewriting, context-dependence correctness, ...
 - e.g. double push-out, Milner's bi-graphs
 - restricting general graph rewriting
 - Synchronised Hyperhedge Replacement (SHR, from Sensoria IP-FP6)
 - Architectural Design Rewriting (ADR, forthcoming)

Implementing concepts in GCM

- when-event-if-cond-then-act list of rules
 - where act either an adaptation or a message to a set of companion managers
 - as JBoss beans



Part II: formalisation





Example: AM ask component f to change location and attach to a new external state (application of 2nd rule - Aldinucci, Tuosto)





Part II: BeSke (orchestration)



ORCHESTRATION OF MANAGEMENT







ORCHESTRATION OF MANAGEMENT

Managers are orchestrated via an overlay network

- in GCM naturally hierarchic (sort of "synch fat-tree")
 - however, the orchestration between children of the same node is not restricted and can be set up according to a user-defined goal
- in general, no restrictions

methodologies to reason about management

- e.g. manager orchestration as service orchestration
 - Orc to describe their orchestration (Misra, Cook, Hoare, ...)
 - reason on Orc programs to prove management global properties
 - semi-formal reasoning for Orc (Aldinucci, Danelutto, Kilpatrick)
 - papers at Europar 07, CoreGRID Symposium 07, IEEE PDP 08, ...





DIFFERENT ORCHESTRATIONS (EXAMPLES)





Part II: Component & Services



COMPONENT, SERVICES OR BOTH?

we re-defined and implemented autonomic BeSke in SCA/Tuscany

- proof-of-concept implementation
- JBoss rule-based manager

few differences

- manager: JBoss rules vs POJO code
- protocols: standard XML/SOAP vs Proactive
- binding: static vs dynamic

proposal for standard extension

- dynamic binding of components
- Tuscany people shown interest







SCA/TUSCANY FARM PERFORMANCES







ANALYSIS: OVERHEADS (GCM/PROACTIVE)





Part II: Static & Dynamic



ANALYSIS: OVERHEAD (ALTERNATIVE IMPL)

ASSIST/C++ overheads (ms)

M. Aldinucci, A. Petrocelli, E. Pistoletti, M. Torquati, M. Vanneschi, L. Veraldi, and C. Zoccolo. Dynamic reconfiguration of grid-aware applications in ASSIST. Euro-Par 2005, vol. 3648 of LNCS, Lisboa, Portugal. Springer Verlag, August 2005.

parmod kind	Data-parallel (with shared state)							Farm (without shared state)					
reconf. kind	add PEs			remove PEs				add PEs			remove PEs		
# of PEs involved	1→2	2→4	4→8	2→1	4→2	8→4		$1 \rightarrow 2$	2→4	4→8	$2 \rightarrow 1$	4→2	8→4
$egin{array}{c} R_l & { m on-barrier} \ R_l & { m on-stream-item} \end{array}$	1.2 4.7	1.6 12.0	2.3 33.9	0.8 3.9	1.4 6.5	3.7 19.1		~ 0	~ 0	~ 0	~ 0	~ 0	~ 0
R_t	24.4	30.5	36.6	21.2	35.3	43.5		24.0	32.7	48.6	17.1	21.6	31.9





IT IS JUST C++ AGAINST JAVA?

No, unfortunately it is not so simple ...

- dynamic class loading (red vs blue zone of the previous chart)
- dynamic introspection
- dynamic binding
- generic data serialisation, shared data alignment
- JIT, code factories, etc.
- non optimised protocols
 - look-ahead resource recruiting
 - pre-deployment
 - atomic multicast (replica management)
 - consensus (reconf-safe-points)
- and at the end ... C++ is usually a bit faster than Java





SUMMING UP ...

exploit both static and dynamic techniques

- represent adaptations as graph transformations
 - in such a way only correct configuration can be generated (e.g. as types)
 - QoS constraints with free variables
- bound free variables with values
 - free variables can be bound at compile, launch time with constant or non constant values
 - manage adaptation accordingly

* uniformly define static and dynamic adaptations

apply them the earlier is possible

compile/deploy/launch/run-time

- here abstraction (e.g. high-level BeSke) become crucial
 - idiom recognition and generative approach





CONCLUSIONS

Behavioural Skeletons

- templates with built-in management for the App designer
- methodology for the skeleton designer
 - management can be changed/refined
 - just prove your own management is correct against skeleton functional description
- can be freely mixed with standard GCM components
- already implemented on GCM (GridCOMP STREP)

#Future work

- many interesting open problems
 - irrespectively of buzzwords (e.g. grid/cyber-infrastructure)
 - irrespectively specific technologies (e.g. component/services)
- this might mean we are trying to address the core of the problems



*** * * ***

THANK YOU

related CoreGRID TR

- M. Aldinucci, M. Danelutto, and P. Kilpatrick. Hierarchical autonomic management: a case study with skeletal systems. CoreGRID Technical Report TR-0127, February 2008.
- 2. P. Kilpatrick, M. Danelutto, M. Aldinucci. Prototyping and reasoning about distributed systems: an Orc based framework. CoreGRID Technical Report TR-0102, August 2007.
- 3. P. Kilpatrick, M. Danelutto, M. Aldinucci. Deriving Grid Applications from Abstract Models. Technical Report TR-0085, April 2007.
- 4. M. Aldinucci, G. Antoniu, M. Danelutto, M. Jan. Fault-tolerant data sharing for high-level grid programming: a hierarchical storage architecture.
 - CoreGRID Technical Report TR-0058, August 2005.
- 5. M, Aldinucci, A. Benoit. Automatic mapping of ASSIST applications using process algebra. CoreGRID Technical Report **TR-0016**, October 2005.
- 6. M. Aldinucci, F. André, J. Buisson, S. Campa, M. Coppola, M. Danelutto, C. Zoccolo. Parallel program/component adaptivity management CoreGRID Technical Report TR-0014, September 2005.
- 7. J. Dünnweber, S. Gorlatch, M. Aldinucci, S. Campa, M. Danelutto. Behavior Customization of Parallel Components for Grid Application Programming. CoreGRID Technical Report TR-0002, April 2005.
- 8. M. Aldinucci, M. Danelutto, J. Dünnweber, S. Gorlatch. Optimization Techniques for Implementing Parallel Skeletons in Distributed Environments.
 - CoreGRID Technical Report TR-0001, January 2005.