Grid programming with components: an advanced COMPonent platform for an effective invisible grid



WP3 - Perspective roadmap for autonomic management

Marco Aldinucci

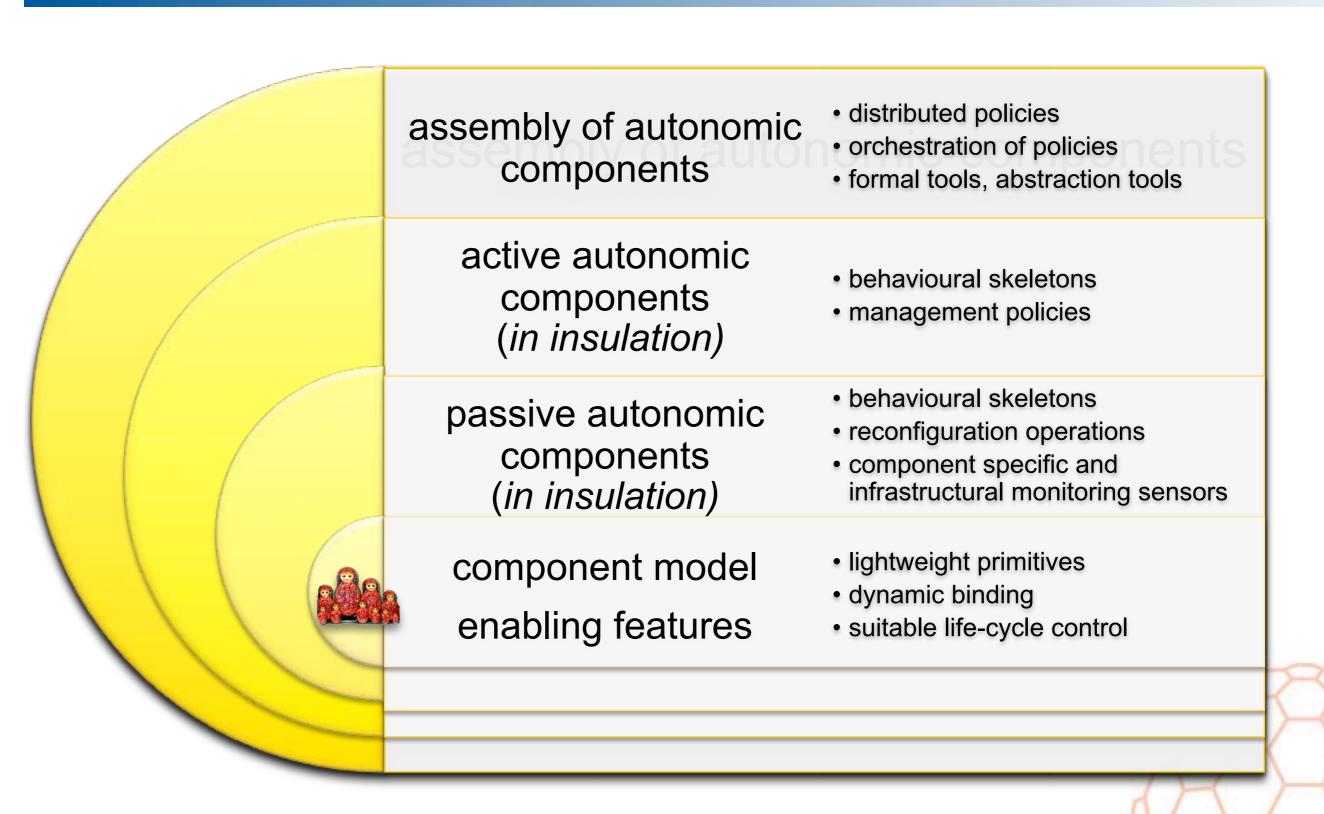
Dept. Computer Science – University of Pisa Programming Model Institute - CoreGRID

aldinuc@di.unipi.it

Unipi & ISTI-CNR working group

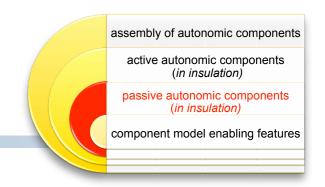
M. Danelutto, D. Laforenza, M. Aldinucci, N. Tonellotto, S. Campa, P. Dazzi, G. Zoppi

Outline





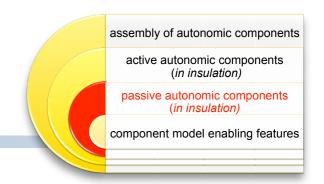
Passive Autonomic Components



- BeSke are parametric composite components
 - extend controllers with ABC & MC
 - non-functional ports for reconfiguration & monitoring
 - implements a modified LC
 - because the manager is a component that cannot be stopped
 - currently provide several flavours of inner components replication
 - stateless farm, stateless farm with initialization, stateless dataparallel, stateful dataparallel
 - enough to cover use cases
- accomplishment state: 95%
 - remaining 5% is mostly related to code reengineering



Farm (functional replication)



- farm BeSke
 - stateless
 - composite with a pair of streaming unicast/from_any ports
 - support composite workers
 - provide parallelism degree change, load-balancing, etc ...
 - autonomic policy for performance assessed
 - ofor performance at least (in insulation)
 - support several use case
 - o e.g. Atos
 - just wrap the sequential code in a primitive component, and farm it out
 - see Nicola's talk & demo

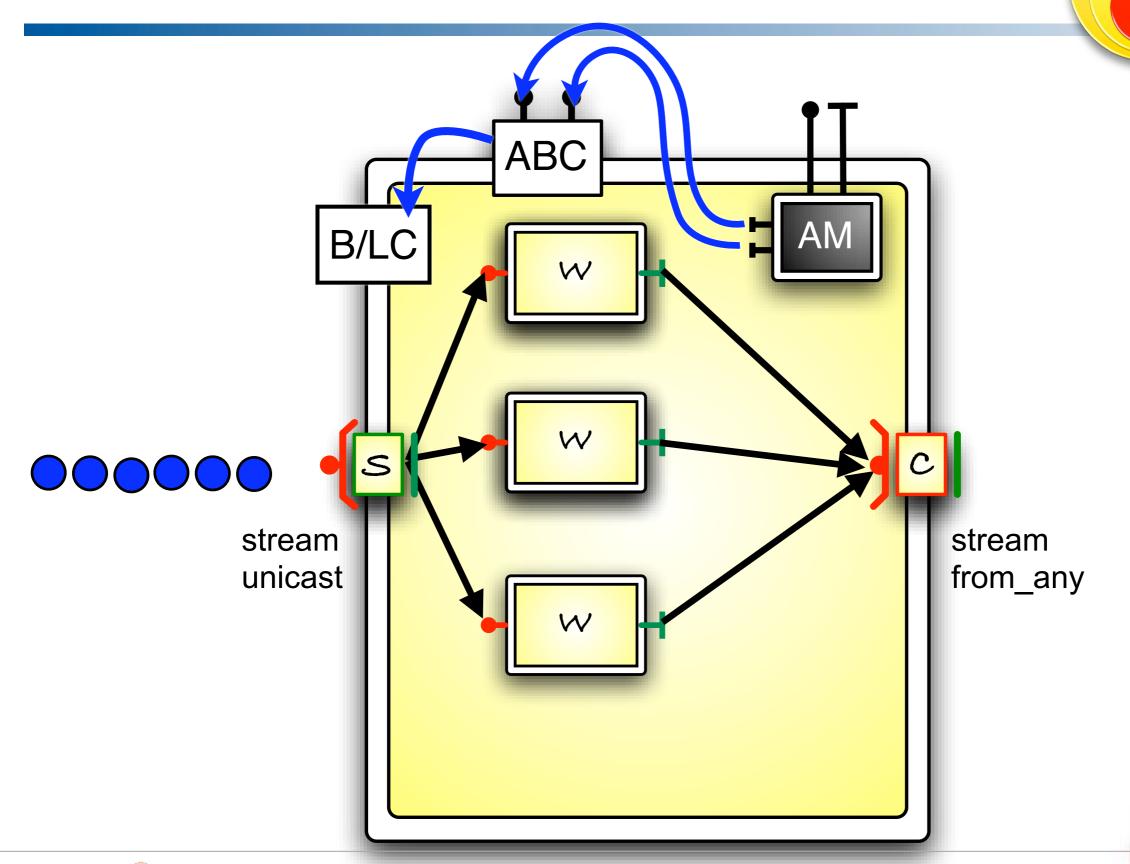


farm (in theory ...)

assembly of autonomic components

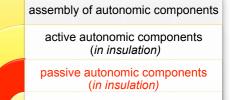
active autonomic components (in insulation)

passive autonomic components (in insulation)

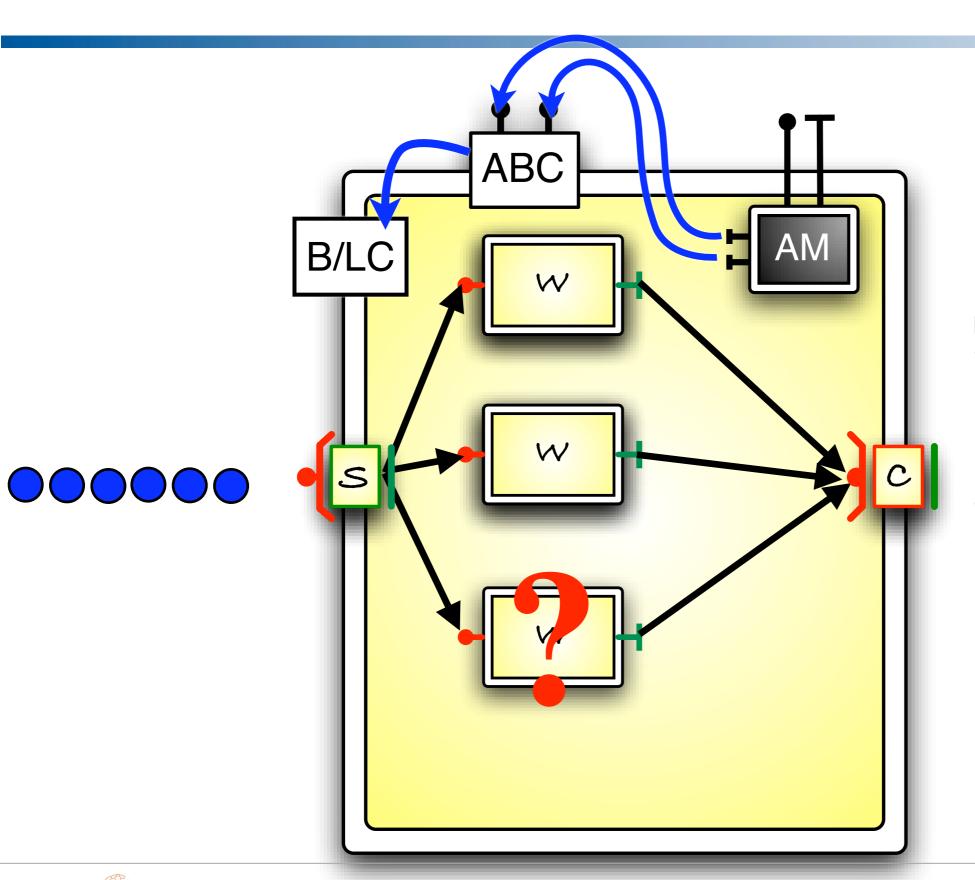




farm (in reality ...)



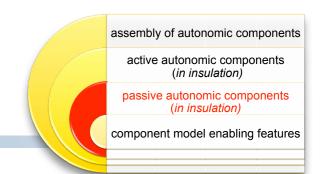
component model enabling features



Notes

- rebalancing operation is logically not needed and it has been introduced because tasks cannot be transiently stopped at the composite membrane
- not sure it works with nonstreaming ports (not void return type)

Data parallel



- data parallel BeSke family
 - stateless and stateful
 - composite exhibiting any number of multicast-gather ports
 - they are bound to inner components according to a fixed predefined parametric pattern
 - e.g. server: scatter client: gather
 - e.g. server: scatter and broadcast client: none
 - autonomic policy for performance assessed
 - for performance at least (in insulation, not yet implemented)
 - supports use case (e.g. IBM)
 - stateful DP provides data redistribution according to user defined&implemented set/get redistribution ports
- o see Sonia's talk & demo

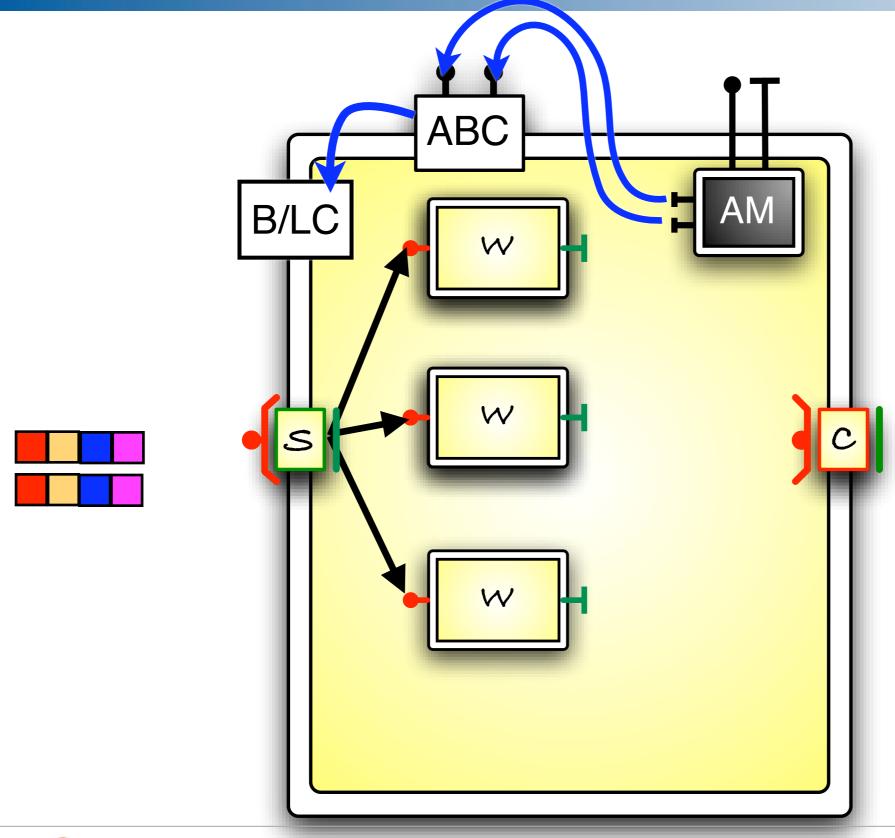


Data Parallel (stateless, pure map)

assembly of autonomic components

active autonomic components (in insulation)

passive autonomic components (in insulation)





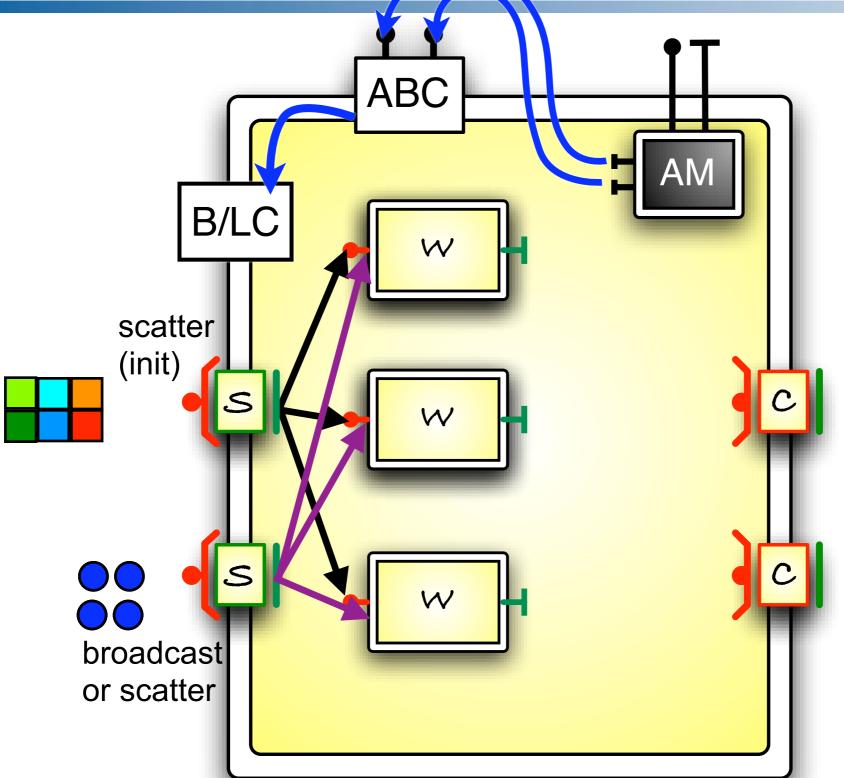
Data Parallel (stateful, distrib state, map)

active autonomic components
(in insulation)

assembly of autonomic components

passive autonomic components (in insulation)

component model enabling features

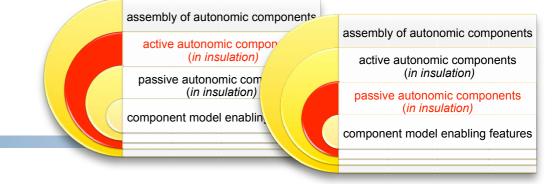


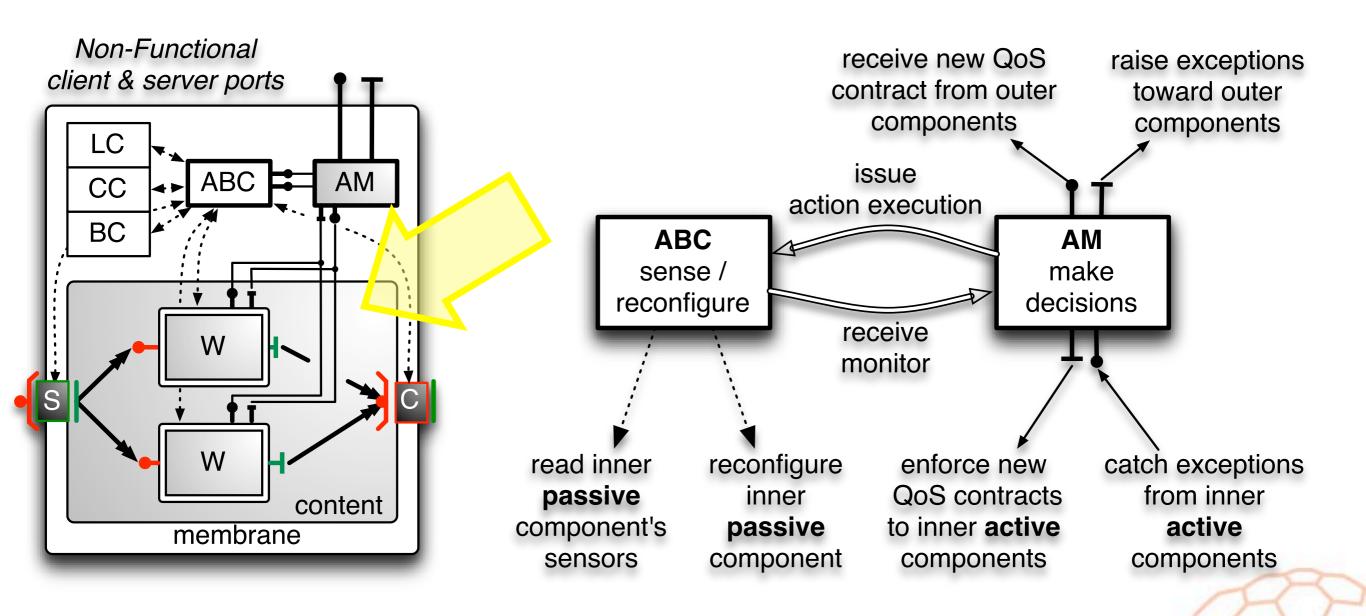
Notes

- any number of server and client ports (either RPC or stream, in theory)
- the model cannot (structurally) enforce init happens before requests on other ports
- port reconfiguration and data redistribution should be atomic (no tasks should be distributed in the middle. We are not sure we can enforce with proactive
- data redistributions are functional requests (they should be because they are related to business logic) but the inner components execute them by way of a NF port in such a way they can be executed in stopped state (workaround)
- enqueued task cannot be cancelled (IBM use case)



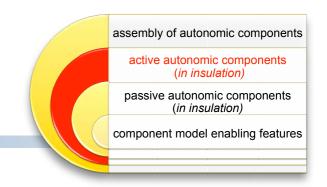
A more defined picture







Active Autonomic Components



- Active = passive + manager
 - non-functional internal and external client and server ports
 - supporting event based orchestration
 - the manager defines a local control policy + co-ordination with other managers
 - when-event-if-cond-then-actions
 - where actions can be invocations to the ABC or messages toward other managers
 - reaction engine can be implemented via JBoss rules
 - already experimented with behavioural skeletons in SCA
 - M. Aldinucci, M. Danelutto, G. Zoppi, P. Kilpatrick.
 Advances in autonomic components & services. Currently submitted



Active management policies - examples

assembly of autonomic components

active autonomic components (in insulation)

passive autonomic components (in insulation)

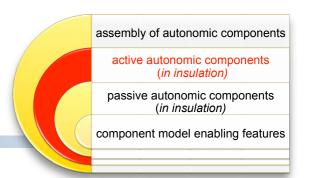
component model enabling features

Component	Manager Contract		m_i	
C_1	active (pipe)		$K_{\text{low}}, K_{\text{high}}$ constants;	$T_{C_2}, T_{C_3}, T_{C_4}$ monitored
	$CP = K_{low} \le T_{self} \le$	K_{high}	$T_{\text{self}} = \max\{T_{C_2}, T_{C_3}, T_{$	
			$\mathcal{CP}_{C_2} = \mathcal{CP}_{C_3} = \mathcal{CP}_{C_4}$	= CP
C_3	active (farm)		IT_{self} = request inter-arriv	al time; n_{self} = #workers
	$CP = (CP_{super}) \wedge (I')$	$T_{\text{self}} \leq T_{\text{self}}$	let C_j children of $C_3, 1 \leq j$	$j \leq n$: T_{C_j} monitored
				$\mathcal{CP}_{C_j} = \text{optimise}(T_{C_j}); \mathcal{G} = 1/n_{\text{self}}$
C_5	active (pipe)		$T_{C_6}, T_{C_7}, T_{C_8}$ monitored	
	$CP = CP_{\text{super}}$		$T_{\text{self}} = \max\{T_{C_6}, T_{C_7}, T_{$	C_{C_8}
			$\mathcal{CP}_6 = \mathcal{CP}_7 = \mathcal{CP}_8 = \text{nu}$	I
$C_{2,4,6,7,8}$	passive (none)		provide $T_{C_{2,4,6,7,8}}$ via NF	port (respectively)
Component	Manager Contract		m_i	
pipe	$\mathcal{CP} = K_{\text{low}} \le T_{\text{self}} \le K_{\text{high}}$		$K_{\text{low}}, K_{\text{high}}$ constants; $\forall C_f \in \text{children(self)}: T_{C_f}$ monitored (service time)	
			$T_{\text{self}} = \max_{\forall C_f} \{T_{C_f}\};$	$\forall C_f : \mathcal{CP}_{C_f} = \mathcal{CP}$
farm	$CP = (CP_{super}) \land (IT_{self} \le T_{self})$		$IT_{ ext{self}}$ monitored (request inter-arrival time); $orall C_f$ \in children(self): T_{C_f} monitored	
				- ,
			tored (service time); $n_{\rm se}$ optimise $(T_{C_f}); \ \mathcal{G} = 1/n_{\rm s}$	$T_{\mathrm{lf}} = \mathrm{\#Workers}; T_{\mathrm{self}} = \sum_{j=1n} T_{C_j}/n; \mathcal{CP}_{C_f} = 0$
Dian		Figure stand Coast	optimise $(T_{C_f}); \ \mathcal{G} = 1/n_{\mathrm{B}}$	elf
Plan		Expected Cost	optimise $(T_{C_f}); \ \mathcal{G} = 1/n_{\mathrm{B}}$	- ,
PL _{F1} move t	he slower worker C_{w} on a platform, if any		optimise $(T_{C_f}); \ \mathcal{G} = 1/n_{\mathrm{B}}$	elf
PL_{F_1} move to faster PL_{F_2} increase		h = cost(stop(optimise $(T_{C_f});\;\mathcal{G}=1/n_{\mathrm{B}}$	Expected Benefit $T_F(\lozenge_{t+h}) = T_{C_x}(\lozenge_t) \triangle$
PL_{F_1} move to faster PL_{F_2} increase cate k PL_{F_3} decrease	olatform, if any se parallelism degree (allo-	h = cost(stop(optimise $(T_{C_f}); \ \mathcal{G} = 1/n_{\mathrm{s}}$ $C_w); \mathrm{deploy}(C_w); \mathrm{start}(C_w))$	Expected Benefit
PLF1 move to faster process cate k places allocate	olatform, if any se parallelism degree (allo- new workers) se parallelism degree (de-	h = cost(stop($\operatorname{optimise}(T_{C_f}); \ \mathcal{G} = 1/n_{\mathrm{s}}$ $C_w); \operatorname{deploy}(C_w); \operatorname{start}(C_w))$ $\operatorname{loy}(C_x); \operatorname{start}(C_w)) \ \text{ for } k \ \text{ in-}$	Expected Benefit $ \text{ decrease service time. } T_F(\lozenge_{t+h}) = T_{C_x}(\lozenge_t) \triangle \\ 0 \leq \Delta \leq 1 \text{ speed difference between the platforms} $
PLF1 move to faster to fas	olatform, if any se parallelism degree (allo- new workers) se parallelism degree (de- e k workers)	h = cost(stop(h = cost(depostances) h = cost(stop(depostances))	$\operatorname{optimise}(T_{C_f}); \ \mathcal{G} = 1/n_{\mathrm{s}}$ $C_w); \operatorname{deploy}(C_w); \operatorname{start}(C_w))$ $\operatorname{loy}(C_x); \operatorname{start}(C_w)) \ \text{ for } k \ \text{ in-}$	Expected Benefit $ \begin{aligned} &\text{decrease service time.} & & & & & & & & & & & & & & & & & & &$
PL_{F_1} move to faster \mathbb{PL}_{F_2} increase cate k PL_{F_3} decrease allocate PL_{F_4} raise VPL_{F_4} move T to a	colatform, if any see parallelism degree (allonew workers) see parallelism degree (dee k workers) iolation stage (C_s) with maximum faster resource, if any see fastest adjacent stages	h = cost(stop(h = cost(depostances) h = cost(stop(depostances)) h = cost(stop(depostances))	$\operatorname{optimise}(T_{C_f}); \ \mathcal{G} = 1/n_s$ $C_w); \operatorname{deploy}(C_w); \operatorname{start}(C_w))$ $\operatorname{loy}(C_x); \operatorname{start}(C_w)) \ \text{for} \ k \ \operatorname{in-}$ $C_w)) \ \text{for} \ k \ \operatorname{in-}$ $C_w)) \ \text{for} \ k \ \operatorname{in-}$	Expected Benefit

M. Aldinucci, M. Danelutto, P. Kilpatrick. Towards hierarchical management of autonomic components: a case study. CoreGRID Technical Report TR-0127, March 2008



Active management - resume



- accomplishment state: 75%
- of the next deliverable
 - equip current demo with a simple manager
 - working in insulation
 - some of those already have in reality
 - code reengineering work mostly
 - farm and data parallel supported
- of the end of the project
 - example of manager coordination
 - probably miming IBM use case
 - dynamic pluggable QoS contracts
 - defined as mobile beans

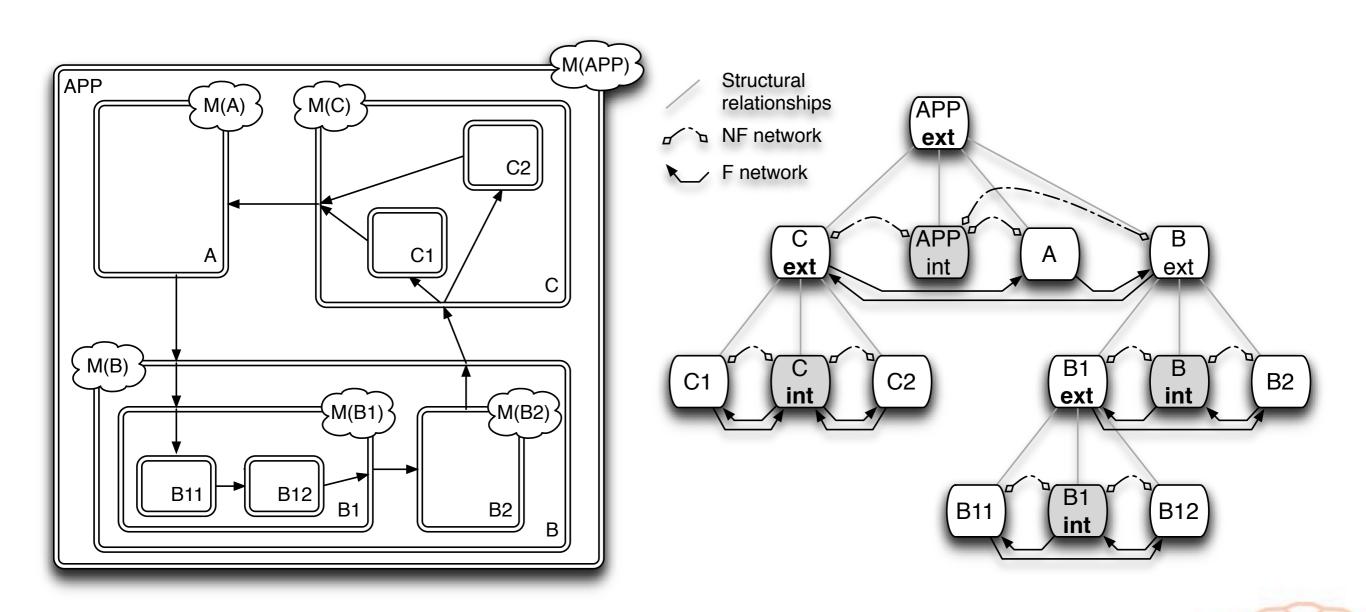


Management orchestration - concepts

assembly of autonomic components

active autonomic components (in insulation)

passive autonomic components (in insulation)





Management orchestration - references

assembly of autonomic components

active autonomic components (in insulation)

passive autonomic components (in insulation)

- accomplishment state: 50%
 - so high just because we don't plan to fully investigate the problem within project lifetime
 - we plan just a single experiment
 - maybe related to IBM use case
 - o in general, we still have many open problems in this regard



GCM Component model thoughts

assembly of autonomic components

active autonomic components
(in insulation)

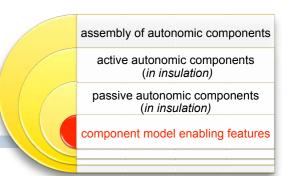
passive autonomic components
(in insulation)

component model enabling features

- possible technical flaws
 - non-uniformity of involved languages
 - ADL, xml, Java, Proactive (or other middleware)
 - lack of static tools
 - correctness of ADL files really error prone
 - factories & automatic generation of non-creative code
 - how many copy-paste do you need to develop an application?
 - o do components actually ease the development?
 - dynamic reservation & deployment support
 - cluster reservation mechanism should be rethought in the light of reconfiguration needs?
 - **G5K**
- conceptual flaws
 - what is the pragmatic of the composite?
 - is reuse really exploited?



GCM/Proactive thoughts & feedbacks



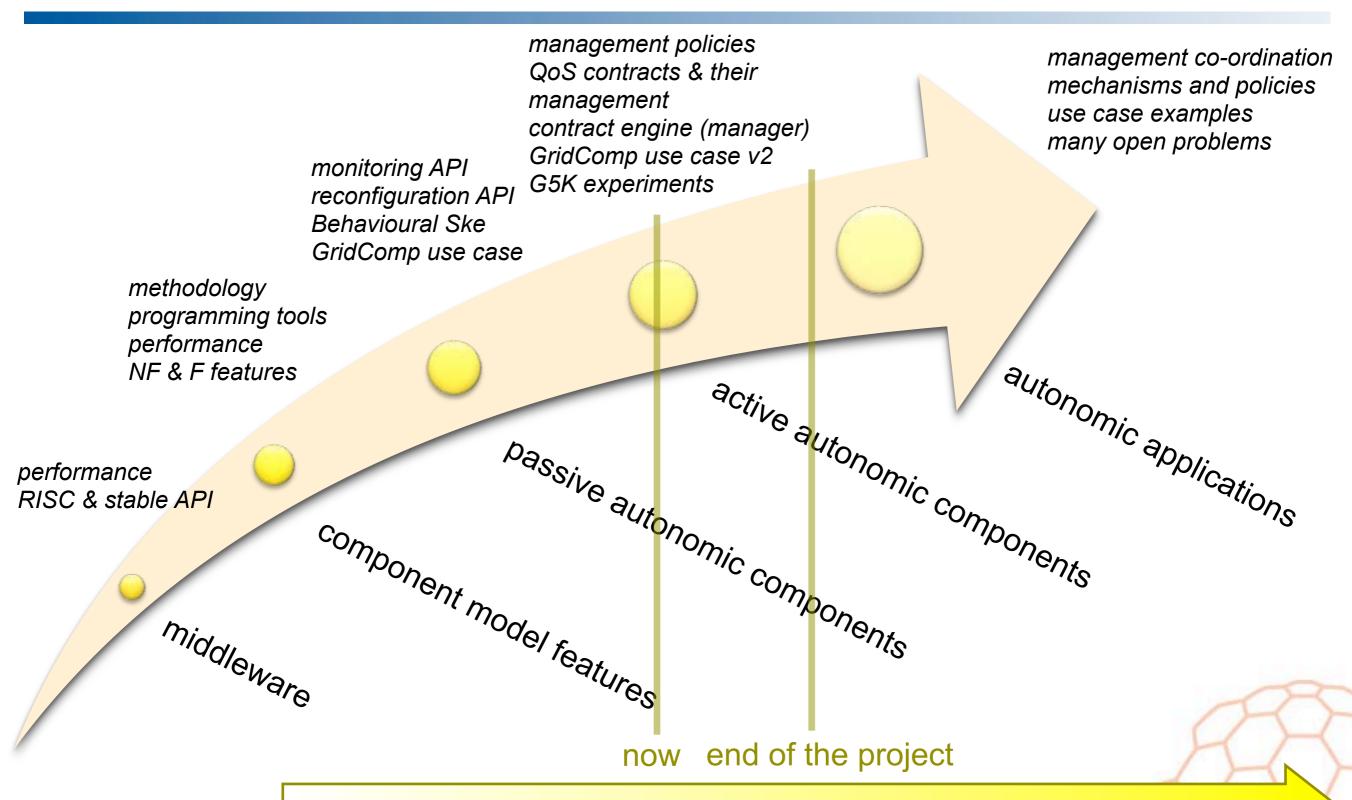
- scarce separation between component model and its run-time support (proactive)
 - the component run-time support is difficult to extend
 - a deep run-time knowledge is strictly required to develop any non-toyish application
- asynchronous communications + queues impair run-time reactivity
 - events (e.g. monitor) refers to current state, but any structural reconfiguration will be applied to non yet enqueued requests
 - e.g. task cancel IBM fingerprint
 - 90% of development effort consists in avoiding unwanted behaviour related to active object model
- life-cycle: stopped/running states probably not enough
 - oproper destruction mechanism, states related to reconfiguration (reconf-safe),
- monitor still missing
 - no real possibility to experiment autonomic policies without it
- GCM has a pretty fat middleware
 - as said in Palma 2007
 - hard to know all the features (then hard to exploit them)
- Proactive versions
 - 3.9_beta
 - any porting to a new version will be complex at this point, apart for expected changes (e.g. monitor)

GCM IMPLEMENTATION STATUS

- **GCM** features under refinement
- *My fat-free (underhanded) wishes
 - * Avoid fat specification
 - Any implementation will hardly be compliant
 - Maybe already too fat
 - Avoid fat implementation
 - Nobody will use it, especially in the HPC community
- Trying to add a "dietetic" QoS control
 - * less possible impact on the middleware, thus if the users don't want it, they should not spend time avoiding it

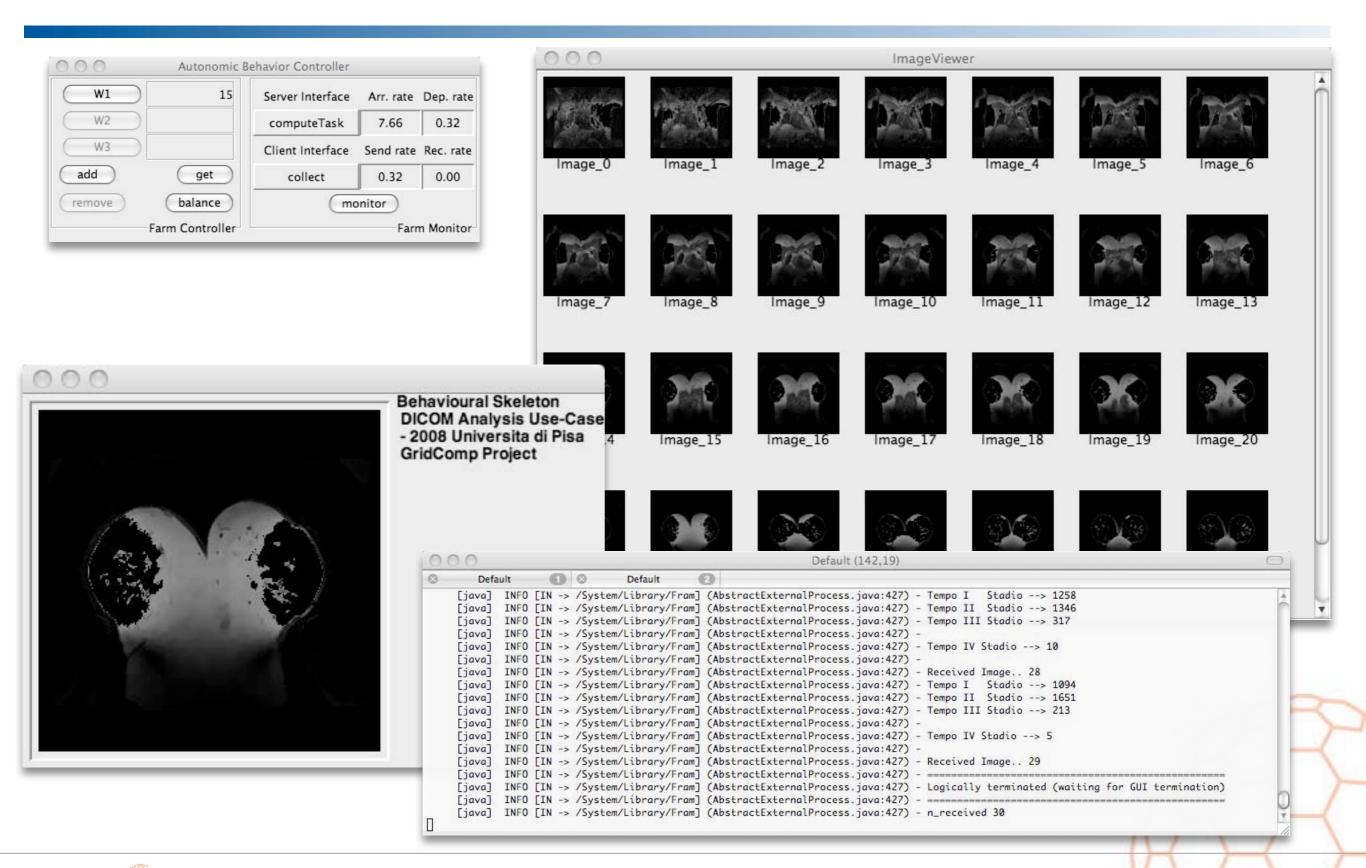


Resume and perspectives





Demo





Resume and perspectives

- passive autonomic components
 - almost all planned features have been implemented
 - LC, ABC, MC, and patches to several controllers
 - behavioural skeletons (functional replication stateless and stateful)
- active autonomic components
 - implementation of management policies for performance (farm and DP)
 - QoS contracts as JBoss rules (static and dynamic)
 - not originally planned, already experimented in SCA
- orchestration of management
 - working on the design
 - formalisation of reconfiguration and QoS contracts
 - formalisation of orchestration strategies
 - uniformly supporting important concepts for HPC and pervasive
 - locality, fault-tolerance, geographic position, ...
- all use cases are supported by BeSke

