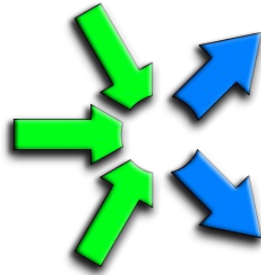


GridKa School 2015

Big Data Virtualization Modern Programming



Introduction to FastFlow programming

Hands-on session

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Outline

- FastFlow basic concepts
 - skeletons and patterns
 - stream concept
 - building blocks
- Core patterns:
 - *pipeline & task-farm*
- High-level patterns:
 - ParallelFor* and Map
 - Macro-DataFlow (mdf)
 - OpenCL StencilReduceLoop
- Targeting distributed system (basic concepts)

The FastFlow tutorial

- Ask for the VM password in case you want to use the Linux VM provided by GridKa
- Update the FastFlow version contained in the VM
 - `cd fastflow`
 - `svn update`
- The FastFlow tutorial is available as pdf file on the GridKa wiki page
 - “FastFlow tutorial”
 - http://wiki.scc.kit.edu/gridkaschool/index.php/Main_Page
- All tests and examples described in the tutorial are available as a separate tarball file: **fftutorial_source_code.tgz**
 - can be downloaded from the wiki page
- In the tutorial source code there are a number of very simple examples covering almost all aspects of using pipeline, farm, ParallelFor, map, mdf, etc..
 - Many features of the FastFlow framework are not covered in the tutorial yet
- There are also a number of small (“more complex“) applications, for example: image filtering, block-based matrix multiplication, mandelbrot set computation, dot-product, etc...



Objectives

- Understand the Structured Parallel Programming methodology
- Have a good idea of the FastFlow framework
 - how it works and its main features
 - also, weakness and strength points
- To be able to write simple FastFlow programs

Structured parallel programming

- Structured parallel programming aims to provide *standard* (and *effective*) rules for composing parallel computations in a machine-independent way
 - **Goal:** *reducing the complexity of parallelization problems by introducing constraints*
 - i.e. restricting the computation structure
 - *Modularity portability* and **programmability** are the keywords
- *Parallel paradigms* are the base components of parallel applications
- Using structured parallel programming force to think parallel
- The programmer is relieved from all concerns related to the implementation, he/she concentrates “*only*” on computational aspects

Separation of concerns principles

Skeletons & Patterns

Algorithmic Skeletons

- From HPC community
- From early '90
- Pre-defined parallel high-order functions provided as constructs or lib calls

- *The same concept at different abstraction levels*
- We use the two terms patterns and skeletons, interchangeably.
 - We want to emphasise the similarities of these two concepts

Parallel Design Patterns

- From SW engineering community
- From early '00
- “Recipes“ to handle parallelism (name, problem, algorithms, solutions, ...)

Using patterns

Available parallel patterns

Map, Reduce, Stencil,....

Pipeline, Task-Farm

Divide&Conquer, MDF,

Data Parallel

Stream Parallel

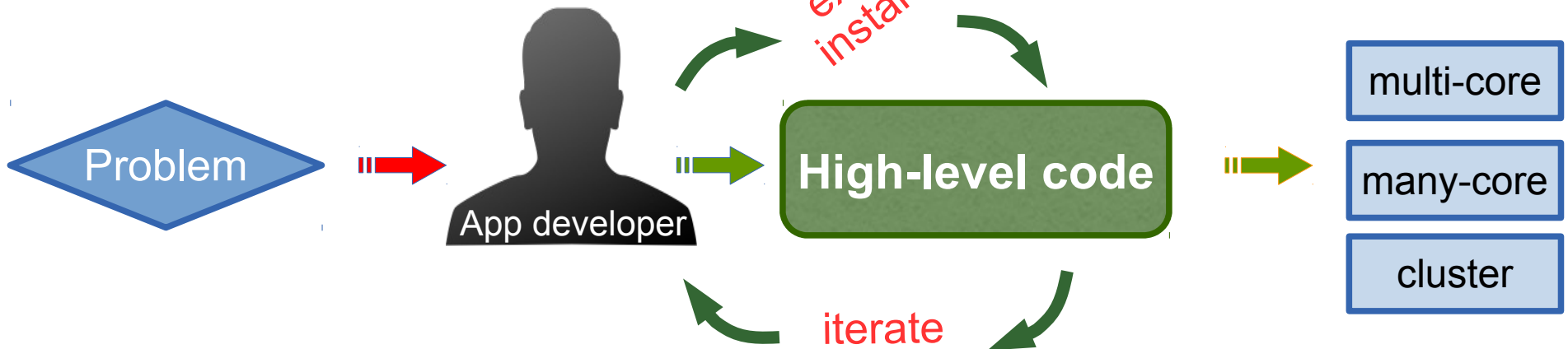
Task Parallel

System developer



run-time support

Which pattern ?



Assessment

Separation of concerns

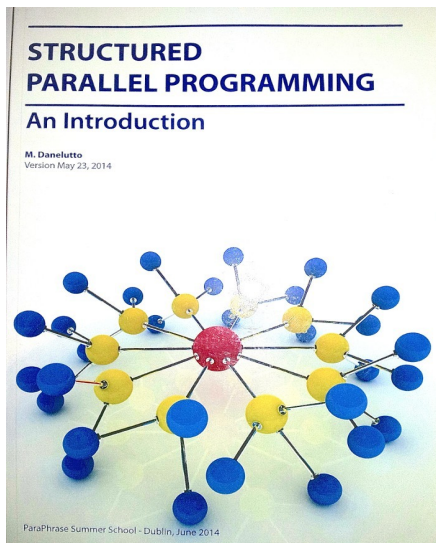
- Application programmer: **what is computed**
- System programmer: **how the result is computed**

Inversion of control

- Program structure suggested by the programmer
- The run-time selects the optimization for the target platform

Performance

- Close to hand tuned code (sometimes better)
- Reduced development time. **Lower total cost to solution.**



“Structured Parallel Programming” by Marco Danelutto

Available on-line as SPM course material at M. Danelutto web page

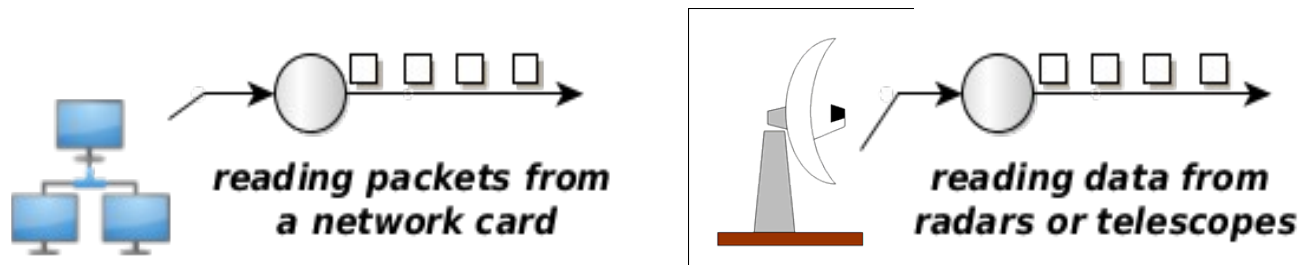
<http://www.di.unipi.it/~marcod>

Stream concept

- The stream concept is important in FastFlow
- A stream is a sequence of values (possibly infinite), coming from a source
 - Stream of images, stream of network packets, stream of matrices, stream of files,
- A streaming application can be seen as a work-flow *graph* whose nodes are computing nodes (sequential or parallel) and arcs are channels bringing streams of data.
- Streams may be either “*primitive*“ (i.e. coming from HW sensors, network interfaces,) or can be generated internally by the application (“*fake stream*”)
- Typically in a stream based computation the first stage receives (or reads) data from a source and produces tasks for next stages.

Stream examples

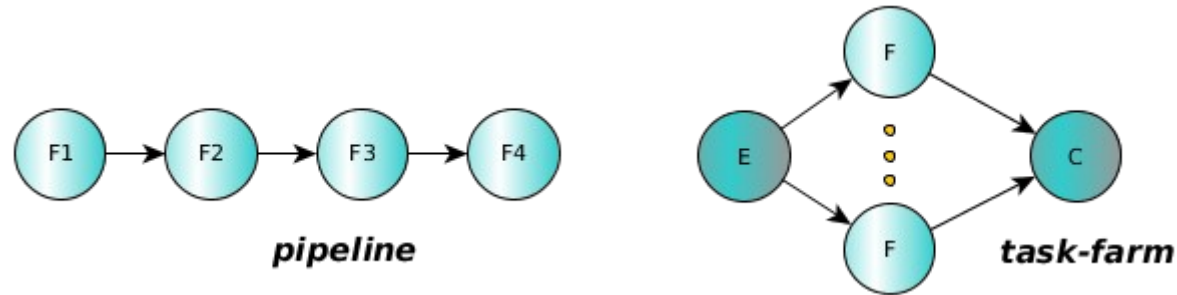
- “*real streams*”



- In these cases it is really important to satisfy minimum processing requirements (bandwidth, latency, etc...) in order to not lose data coming from the source
- “*fake streams*”: streams produced by unrolling loops
 - You don't have an “infinite” source of data
 - The source is a software module

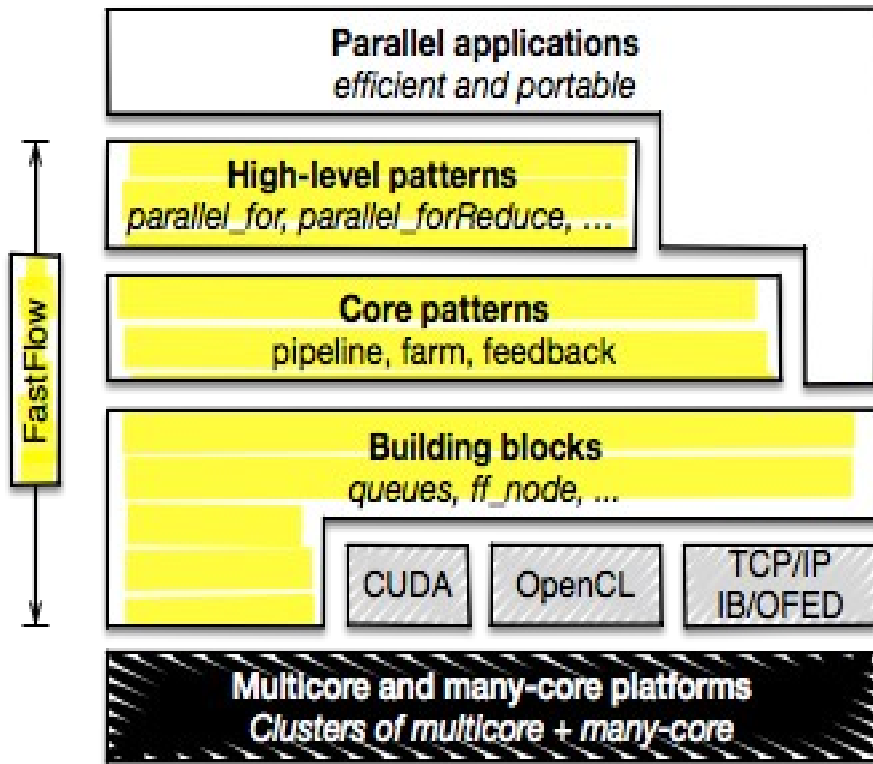
```
for(i=start; i<stop; i+=step)
  allocate data for a task
  create a task
  send out the task
```

Patterns operating on stream



- **pipeline**: computes $F_4(F_3(F_2(F_1(x))))$ for each x
 - Pipeline computing elements are called *stages*
- **task-farm** (or *farm*), **models functional replication**
 - Sometimes called also “master-worker”
 - Computing elements called: Emitter (E), Worker (computing F) and Collector (C)
 - The Emitter, schedules tasks towards the Workers
 - The Collector, gathers tasks from Workers

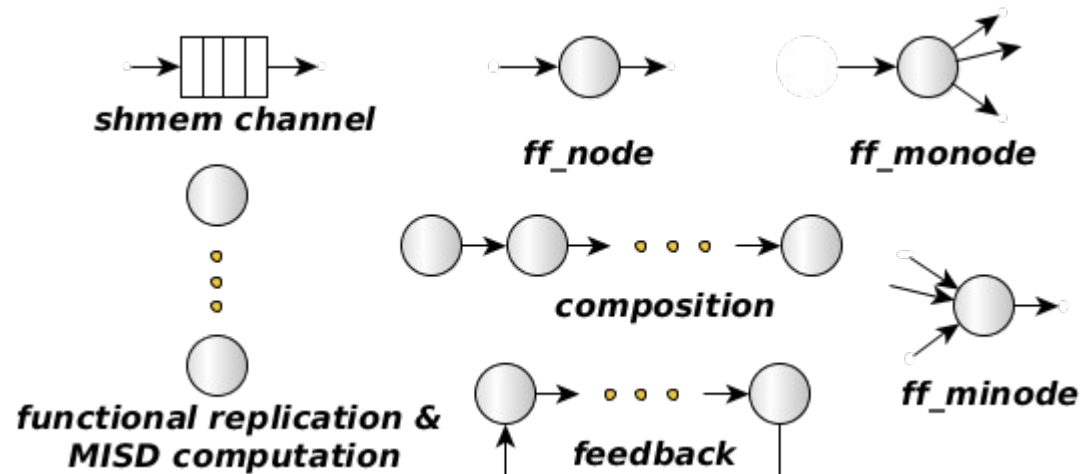
The FastFlow layers



<http://mc-fastflow.sourceforge.net>
<http://calvados.di.unipi.it/fastflow>

- C++ class library
- Promotes (high-level) structured parallel programming
- It aims to be flexible and efficient enough to target **multi-core, many-core** and **distributed heterogeneous systems**.
- Layered design:
 - **Building blocks** minimal set of mechanisms: channels, code wrappers, combinators.
 - **Core patterns** streaming patterns (*pipeline* and *task-farm*) plus the *feedback* pattern modifier
 - **High-level patterns** aim to provide flexible reusable parametric patterns for solving specific parallel problems

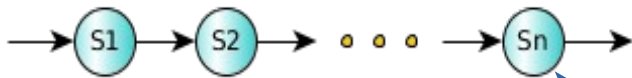
Building blocks



- Minimal set of efficient mechanisms and functionalities
- Nodes are concurrent entities (i.e. POSIX threads)
- Arrows are channels implemented as SPSC lock-free queue
 - bounded or unbounded in size

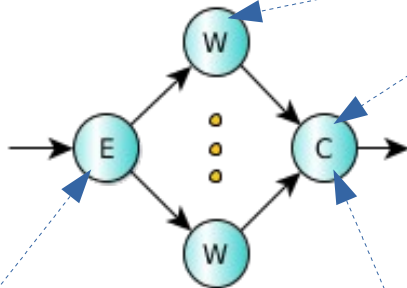
Stream Parallel Patterns (“core” patterns)

pipeline



```
ff_Pipe<myTask> pipe(S1,S2,...,Sn);  
pipe.run_and_wait_end();
```

task-farm



ff_node

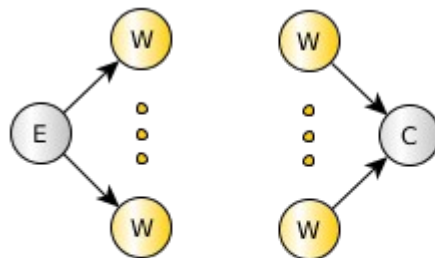
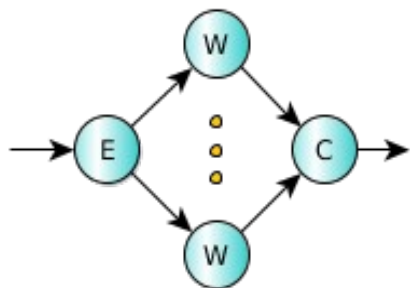
```
std::vector<std::unique_ptr<ff_node> > W;  
myEmitter E;  
myCollector C;  
  
ff_Farm<myTask> farm(std::move(W), E, C);  
farm.run_and_wait_end();
```

Emitter:
schedules input data items

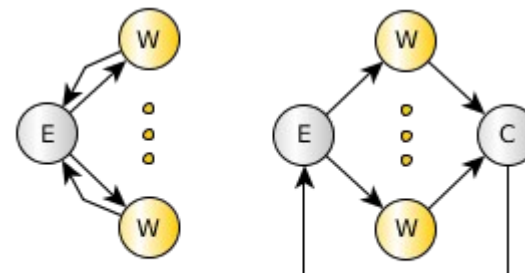
Collector:
gathers results

Stream Parallel Patterns (“core” patterns)

task-farm

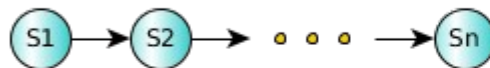
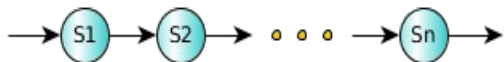


task-farm

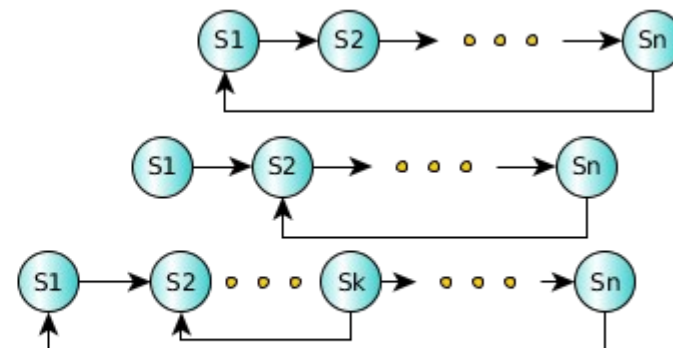


task-farm + feedback

pipeline



pipeline

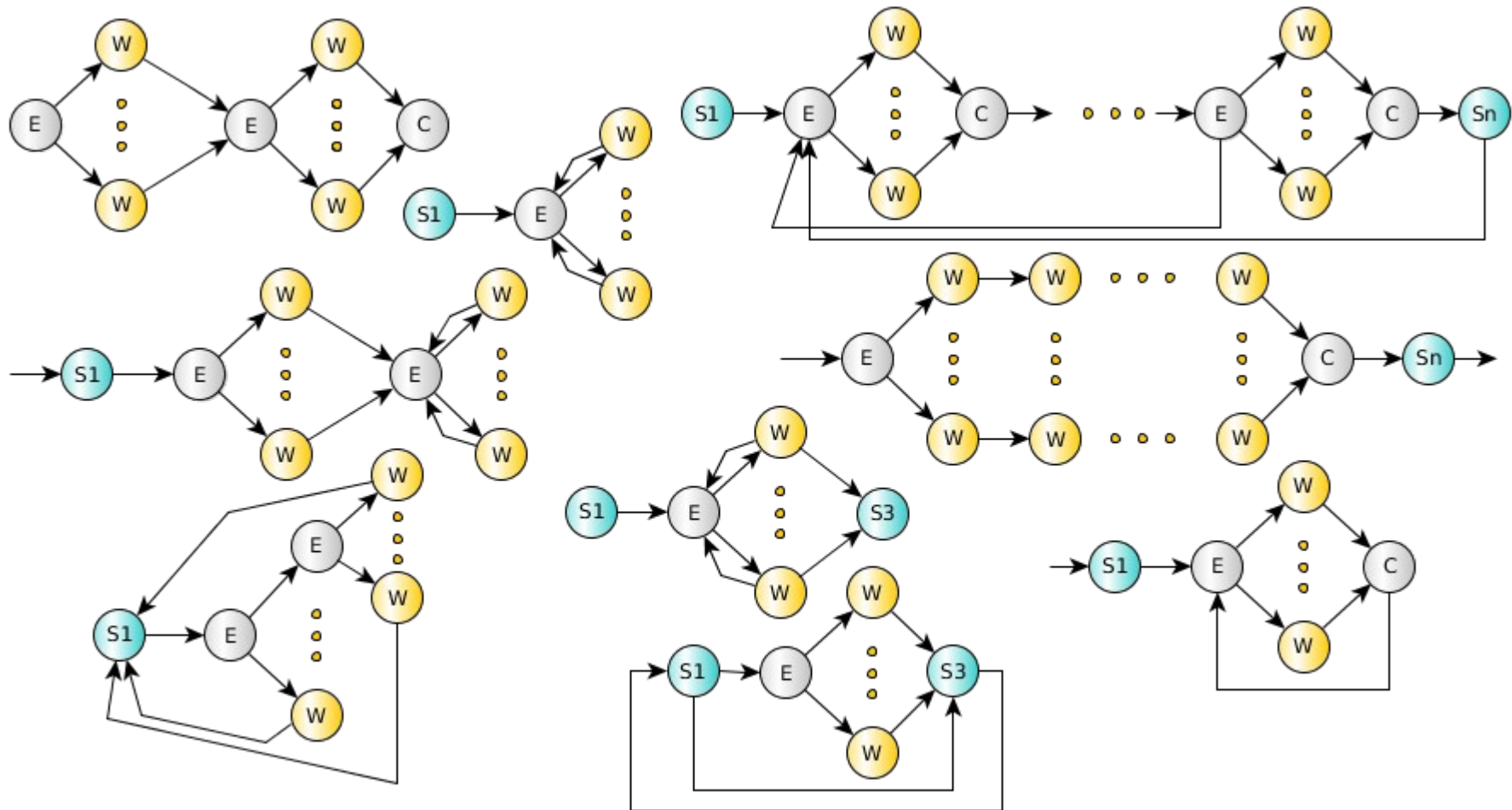


pipeline + feedback

Specializations

Patterns

Core patterns composition



pipeline + task-farm + feedback

Core patterns: sequential *ff_node*

code wrapper pattern

```
struct myNode: ff_node_t<TIN,TOUT> {
  int svc_init() { // optional
    // called once for initialization purposes
    return 0; // <0 means error
  }
  TOUT *svc(TIN * task) {
    // do something on the input task
    // called each time a task is available
    return task; // also EOS, GO_ON, ....
  };
  void svc_end() {
    // called once for termination purposes
    // called if EOS is either received in input
    // or it is generated by the node
  }
};
```

- A sequential *ff_node* is an active object (thread)
- Input/Output tasks (stream elements) are memory pointers
- The user is responsible for memory allocation/deallocation of tasks
 - FF provides a memory allocator (not introduced here)
- Special return values:
 - *EOS* means End-Of-Stream
 - *GO_ON* means “I have no more tasks to send out, give me another input task (if any)”

ff_node: generating and absorbing tasks

code wrapper pattern

```
struct myNode1: ff_node_t<Task> {  
    Task *svc(Task *) {  
        // generates N tasks and then EOS  
        for(long i=0;i<N; ++i)  
            ff_send_out(new Task);  
        return EOS;  
    };  
};
```

```
struct myNode2: ff_node_t<Task> {  
    Task *svc(Task * task) {  
        // do something with the task  
        do_Work(task);  
        delete task;  
        return GO_ON; // it does not send out task  
    };  
};
```

- Typically myNode1 is the first stage of a pipeline, it produces tasks by using the *ff_send_out* method or simply returning task from the svc method
- Typically myNode2 is the last stage of a pipeline computation, it gets in input tasks without producing any outputs

Core patterns: *ff_pipe*

pipeline pattern

```
struct myNode1: ff_node_t<myTask> {
    myTask *svc(myTask *) {
        for(long i=0;i<10;++i)
            ff_send_out(new myTask(i));
        return EOS;
    }
};
struct myNode2: ff_node_t<myTask> {
    myTask *svc(void *task) {
        return task;
    }
};
struct myNode3: ff_node_t<myTask> {
    myTask *svc(void * task) {
        f3(task);
        return GO_ON;
    }
};
myNode1 _1;
myNode2 _2;
myNode3 _3;
ff_Pipe<> pipe(_1,_2,_3);
pipe.run_and_wait_end();
```

- *pipeline* stages are *ff_node(s)*
- A *pipeline* itself is an *ff_node*
 - It is easy to build pipe of pipe
- **ff_send_out** can be used to generate a stream of tasks
- Here, the first stage generates 10 tasks and then EOS
- The second stage just produces in output the received task
- Finally, the third stage applies the function f3 to each stream element and does not return any tasks



Simple *ff_pipe* examples

- Let's comment on the code of the 2 simple tests presented in the FastFlow tutorial:
 - `hello_pipe.cpp`
 - `hello_pipe2.cpp`

Core patterns: *ff_farm*

(1)

task-farm pattern

```
struct myNode: ff_node_t<myTask> {  
    myTask *svc(myTask * t) {  
        F(t);  
        return GO_ON;  
    }  
};
```

```
std::vector<std::unique_ptr<ff_node>> W;  
W.push_back(make_unique<myNode>());  
W.push_back(make_unique<myNode>());
```

```
ff_Farm<myTask>  
    myFarm(std::move(W));
```

```
ff_Pipe<myTask>  
    pipe(_1, myFarm, <...other stages...>);  
  
pipe.run_and_wait_end();
```

- Farm's workers are `ff_node(s)` provided via an `std::vector`
- By providing different `ff_node(s)` it is easy to build a MISD farm (each worker computes a different function)
- By default the farm has an Emitter and a Collector, the Collector can be removed using:
 - `myFarm.remove_collector();`
- Emitter and Collector may be redefined by providing suitable `ff_node` objects
- Default task scheduling is pseudo round-robin
- Auto-scheduling:
 - `myFarm.set_scheduling_ondemand();`
- Possibility to implement user's specific scheduling strategies (`ff_send_out_to`)
- Farms and pipelines can be nested and composed in any way

Core patterns: *ff_farm*

(2)

task-farm pattern

```
myTask *F(myTask * t,ff_node*const) {  
  .... <work on t> ....  
  return t;  
}
```

```
ff_Farm<myTask> myFarm(F, 5);
```

```
myTask *F(myTask * t,ff_node*const) {  
  .... <work on t> ....  
  return t;  
}
```

```
ff_OFarm<myTask> myFarm(F, 5);
```

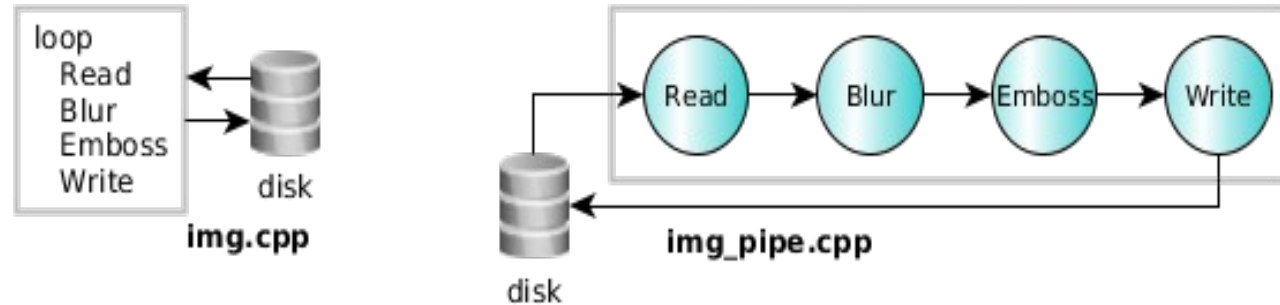
- Simpler syntax
- By providing a function having a suitable signature together with the number of replicas
 - 5 replicas in the code aside
- Default scheduling or auto-scheduling

- Ordered task-farm pattern
- Tasks are produced in output in the same order as they arrive in input
- In this case it is not possible to redefine the scheduling policy

Simple *ff_farm* examples

- Let's comment on the code of the 2 simple tests presented in the FastFlow tutorial:
 - `hello_farm.cpp`
 - `hello_farm2.cpp`
- Then, let's take a look at how to define Emitter and Collector in a farm:
 - `hello_farm3.cpp`
- A farm in a pipeline without the Collector:
 - `hello_farm4.cpp`

Examples: image filtering (img.cpp & img_pipe.cpp)



```
// 4-stage pipeline  
ff_Pipe<Task> pipe( read, blur, emboss, write );  
pipe.run_and_wait_end();
```

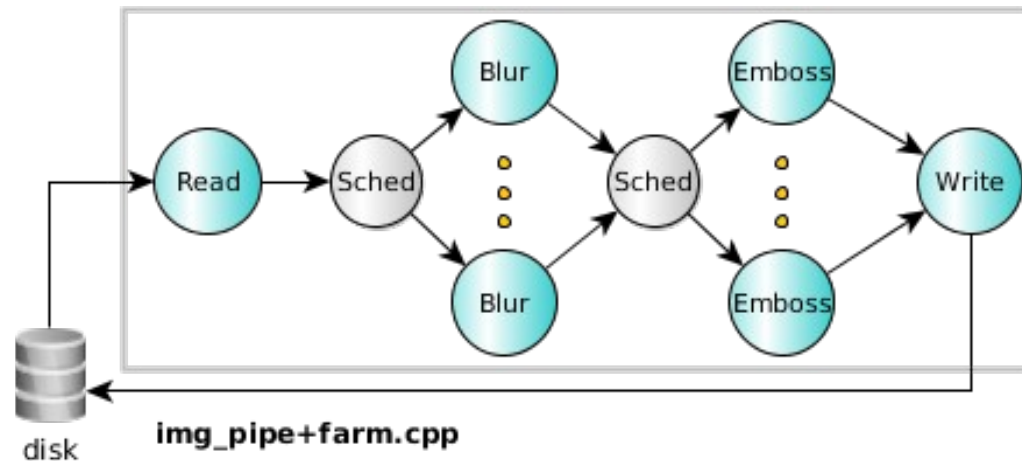
```
// 1st stage  
struct Read: ff_node_t<Task> {  
    Task *svc(Task *) {  
        for(long i=0;i<num_images;++)  
            Image *img = new Image;  
            Img->read(filename);  
            Task *task = new Task(img,filename);  
            ff_send_out(task);  
    }  
    return EOS; // End-Of-Stream  
};
```

```
// 2nd stage  
Task *BlurFilter(Task *in, ff_node*const) {  
    in->image->blur(); return in;  
}
```

```
// 3rd stage  
Task *EmbossFilter(Task *in, ff_node*const) {  
    in->image->blur(); return in;  
}
```

```
// 4th stage  
Task *Write(Task *in, ff_node*const) {  
    in->image->write(in->name);  
    delete in->image;  
    delete in;  
    return reinterpret_cast<Task*>(GO_ON);  
}
```


Examples: image filtering (img_pipe+farm.cpp)



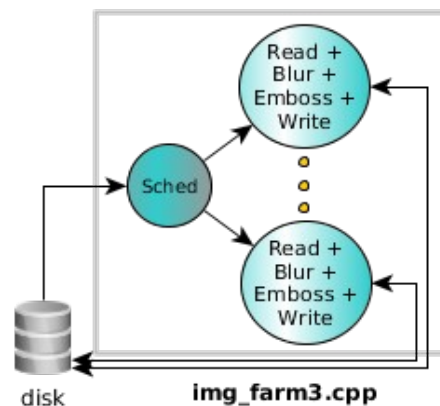
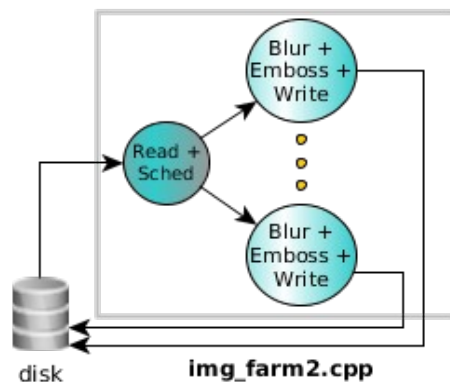
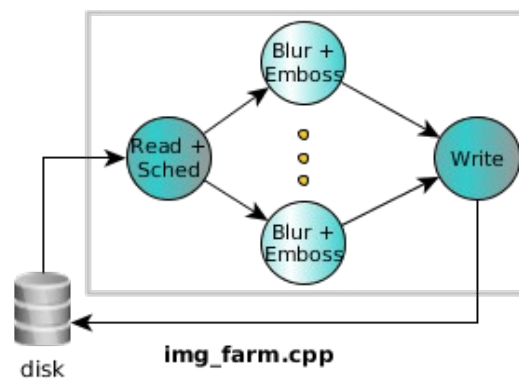
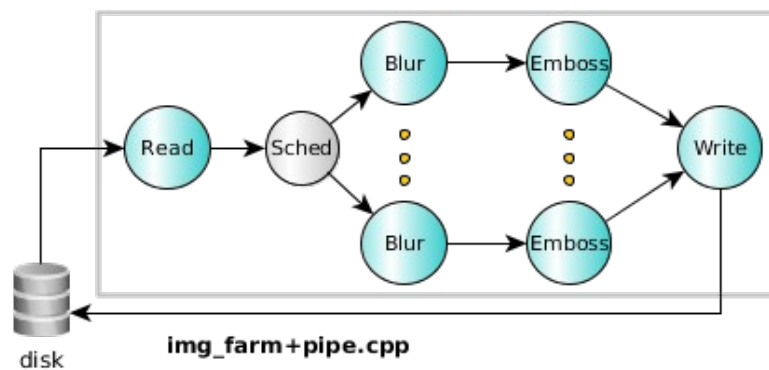
```
ff_Farm<Task> farmBlur(BlurFilter, numBlurWorkers);  
farmBlur.remove_collector();  
ff_Farm<Task> farmEmboss(EmbossFilter, numEmbosWorkers);  
// 4-stage pipeline  
ff_Pipe<Task> pipe( read, farmBlur, farmEmboss, write );  
pipe.run_and_wait_end();
```

```
// ff_node wrapper to the Write function  
struct Writer: ff_minode_t<Task> {  
    Task *svc(Task *task) {  
        return Write(task, this);  
    };  
};
```

Other nodes are the same as before

Examples: image filtering

Other simple transformations



Take a look at the code and how it works !

Proposed exercises using ff_pipe & ff_farm

- Simple file compressor using miniz.c:
 - The sequential implementation is given *simplecomp.cpp*
 - The task is to first implement a pipeline version (read → compute → write), then transform the sequential compute stage in a task-farm.
 - `simplecomp_pipe.cpp`
 - `simplecomp_farm.cpp`
 - **HINT:** the structure is quite similar to the one used in the `img_pipe.cpp` and `img_farm.cpp`, respectively.
- A more complex and efficient implementation is left as homework
 - We will discuss together some possible solutions

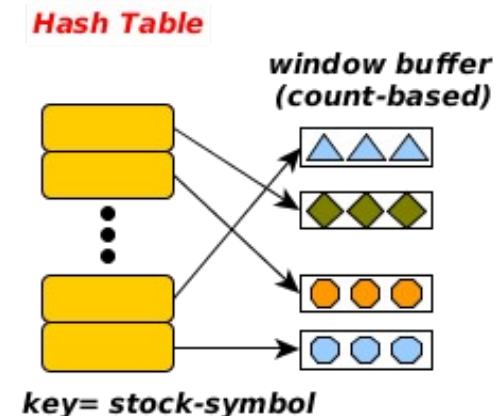
Example: statefull pipeline

- Simplified financial application



Sequential pseudo-code:

```
Receiver rec(port); // recv quotes from the market
while( recv.receive(quote) ) {
  filterQuote(quote); // filters data
  If ((wid=winManagement(quote, win_size, win_slide)) != -1) {
    computeWindow(wid, result); // data ready
    writeOnDisk(result); // write result
  }
}
```



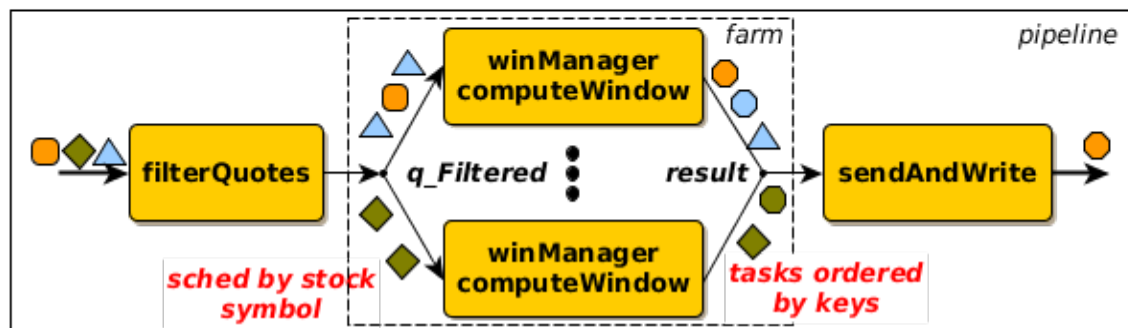
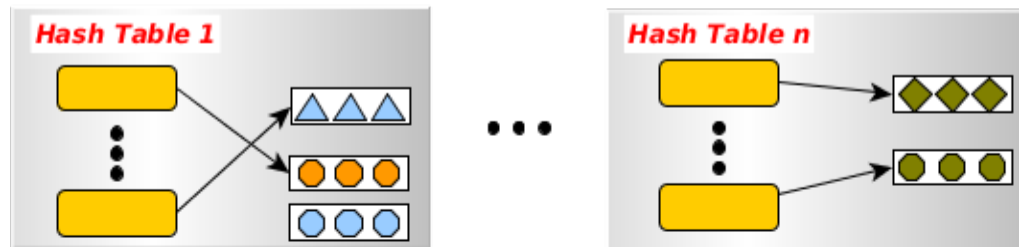
- WindowManagement is based on a hash-table containing different buffers for each stock symbol
- The application is logically a 3-stage pipeline (receive, compute, write) but the middle stage cannot be replicated unless

Example: statefull pipeline

- ... unless ... **the hast-table is *partitioned*** among all workers and the quotes are scheduled by stock symbol

Parallel structure:

```
struct firstStage: ff_node_t<quote_t> {  
    quote_t *svc(quote_t *in) { return filter(*in); }  
};  
....  
Receiver rec(port);  
firstStage first(rec);  
  
std::vector<std::unique_ptr<ff_node>> W;  
for(long i=0;i<nworkers;++i)  
    W.push_back(make_unique<compute>  
                (win_size,win_slide));  
ff_Farm<task_t,ret_t> farm(std::move(W));  
Scheduler<decltype<SchedF> Sched(schedF);  
farm.add_emitter(Sched);  
farm.remove_collector();  
  
lastStage last(writerOnDisk);  
  
ff_Pipe<> pipe(first, farm, last);  
pipe.run_and_wait_end();
```



Potential load balancing problems !

High-level patterns

- Here we consider
 - *ParallelFor** data-parallel patterns
 - *ff_Map*, that can be used in pipeline and task-farm
 - *Macro-Data-flow* (MDF) as data-flow pattern (or task-parallel pattern)
- *Pipeline* and *task-farm* are high-level patterns as well !

- Other patterns available in FastFlow are:
 - *PoolEvolution* for modelling evolutionary applications
 - *StencilReduceLoop* patterns for iterative stencil-like computation (multi-core, CUDA/OpenCL-based GPUs)
 - We will introduce some basic concept later

High-level patterns: ParallelFor

map pattern

```
// sequential code
for(long i=0;i<N; i+=2)
    A[i] = f(i);
```

```
// parallel code
ParallelFor pf;
pf.parallel_for(0, N, 2, [&A](const long i) {
    A[i] = f(i);
});
```

- Loops with independent iterations may be parallelised using the ParallelFor pattern
- The ParallelFor interface is in the *parallel_for.hpp* file
- It is implemented on top of the task-farm with a suitable scheduling strategy
- There are many different methods that can be used
- Iteration scheduling provided:
 - Default static scheduling
 - Static scheduling with interleaving by using **parallel_for_static**
 - Dynamic scheduling
- Also provides *active scheduling* (by using farm's Emitter) and *passive scheduling*

High-level patterns: ParallelForReduce

map-reduce pattern

```
// sequential code: summing all elements  
// of an array  
double sum=0.0;  
for(long i=0;i<N; i++)  
    sum += A[i];
```

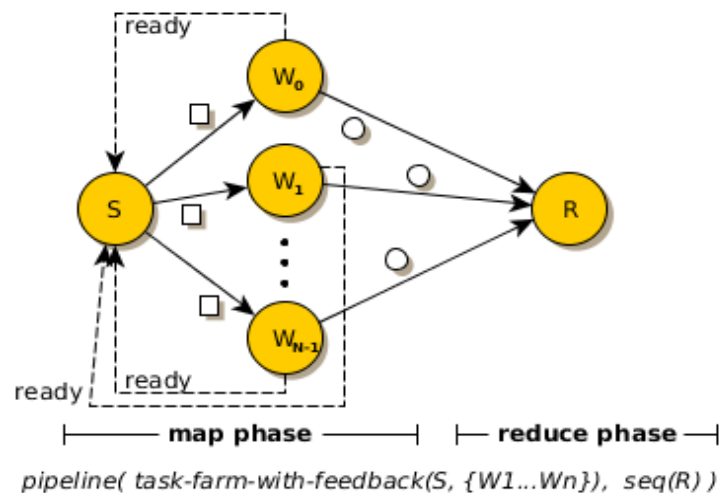
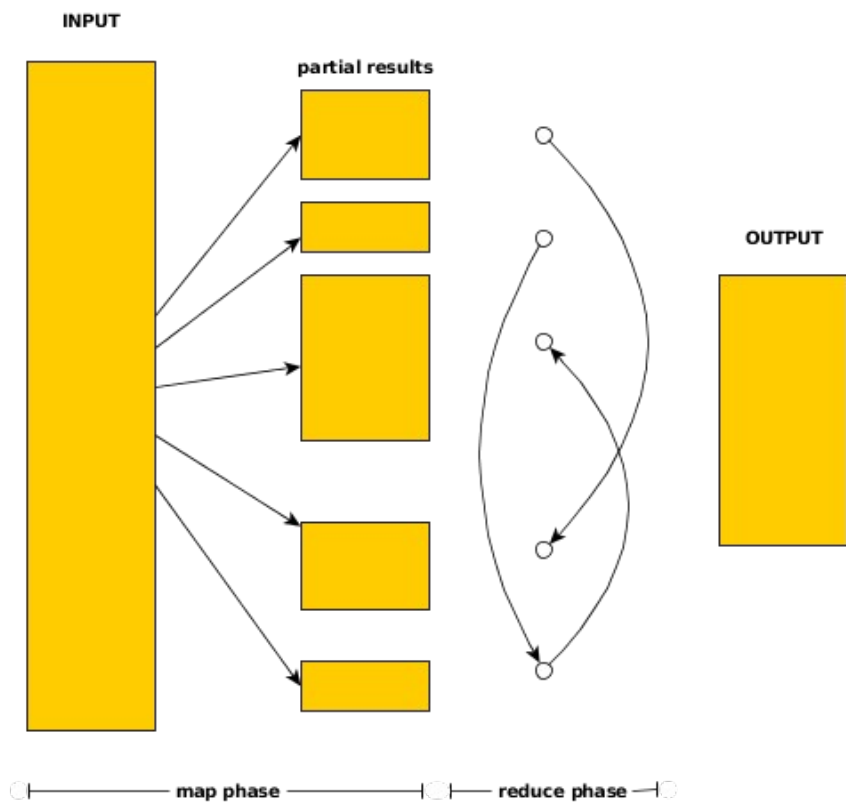
```
// parallel code  
ParallelForReduce<double> pfr;  
pfr.parallel_for_reduce(sum, 0.0, 0,N,  
    [&A](const long i, double &sum) {  
        sum +=A[i];  
    },  
    [](double &sum, const double v) {  
        sum+=v;  
    }  
);
```

- A ParallelFor plus a reduction operation
 - associative operation
- The ParallelForReduce interface is in the *parallel_for.hpp* file
- It is implemented on top of the task-farm with a suitable scheduling strategy
- Executes a local reduction in the body part using a private variable plus a final reduction operation using a combination function.
- Scheduling strategies are the same as those provided by the ParallelFor pattern

High-level patterns: ParallelForPipeReduce

map-reduce pattern

- Useful when the reduction function has to be executed serially
- Offers the option to compute the Map part in pipeline with the Reduce part.
- The ParallelForPipeReduce interface is in the *parallel_for.hpp* file
- Scheduling strategies are the same as those provided by the ParallelFor pattern



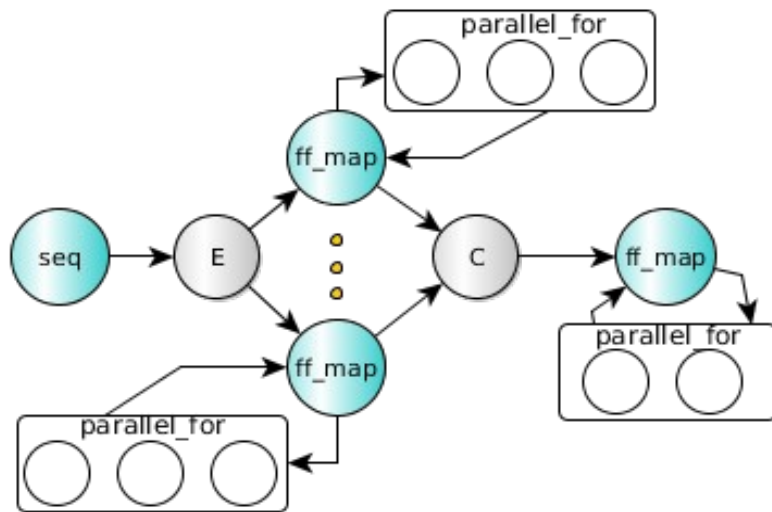


Simple tests using a ParallelFor

- Let's comment on the code of the 2 simple tests presented in the FastFlow tutorial:
 - `hello_parfor.cpp`
 - `arraysum.cpp`

High-level patterns: ff_Map

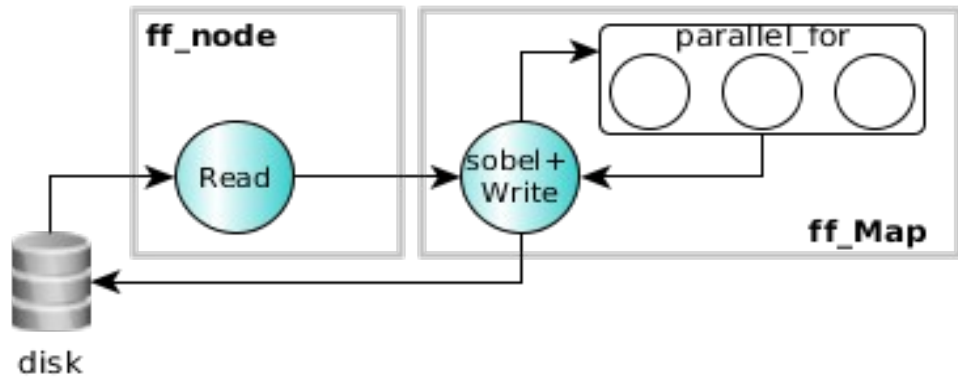
map pattern inside
stream parallel patterns



- The *ff_Map* is just an *ff_node* that wraps a ParallelForReduce pattern
- The *ff_Map* can be used as a pipeline stage and as a farm worker
- It is better to use the *ff_Map* than a plain ParallelFor in a pipeline or farm computations because the runtime knows that the given stage/worker is parallel
 - Better thread mapping strategies and optimizations can be applied

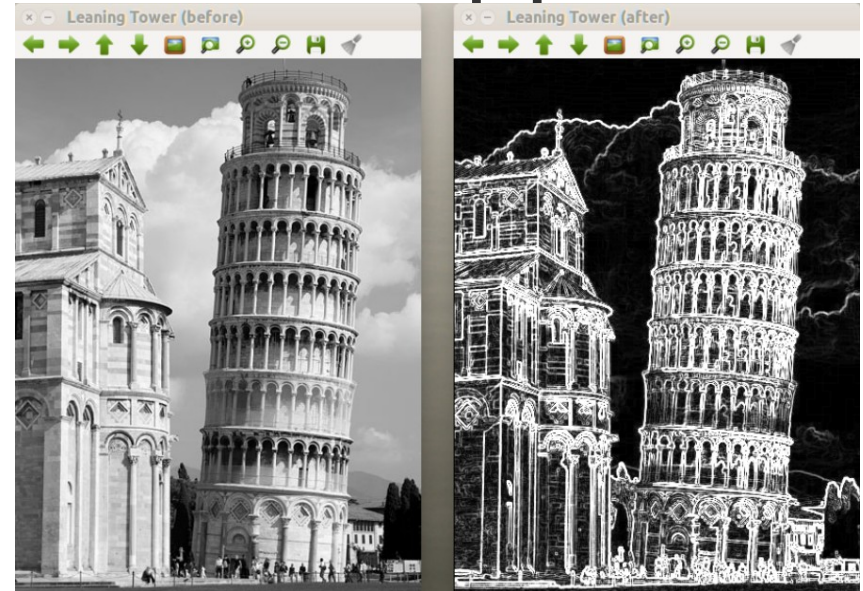
```
struct mapWorker: ff_Map<> {
    void *svc(void *task) {
        ....
        ff_Map<>::parallel_for(...);
        ff_Map<>::parallel_reduce(...);
        return task;
    }
};
```

Examples: Sobel filter (ffsobel.cpp)



```
struct sobelStage: ff_Map<Task> {
    sobelStage(int mapwrks):
        ff_Map<Task>(mapwrks, true) {};

    Task *svc(Task*task) {
        Mat src = *task->src, dst= *task->dst;
        ff_Map<>::parallel_for(1,src,src.row-1,
            [src,&dst](const long y) {
                for(long x=1;x<src.cols-1;++x) {
                    .....
                    dst.at<x,y> = sum;
                }
            });
        const std::string outfile="./out"+task->name;
        imwrite(outfile, dst);
    }
}
```

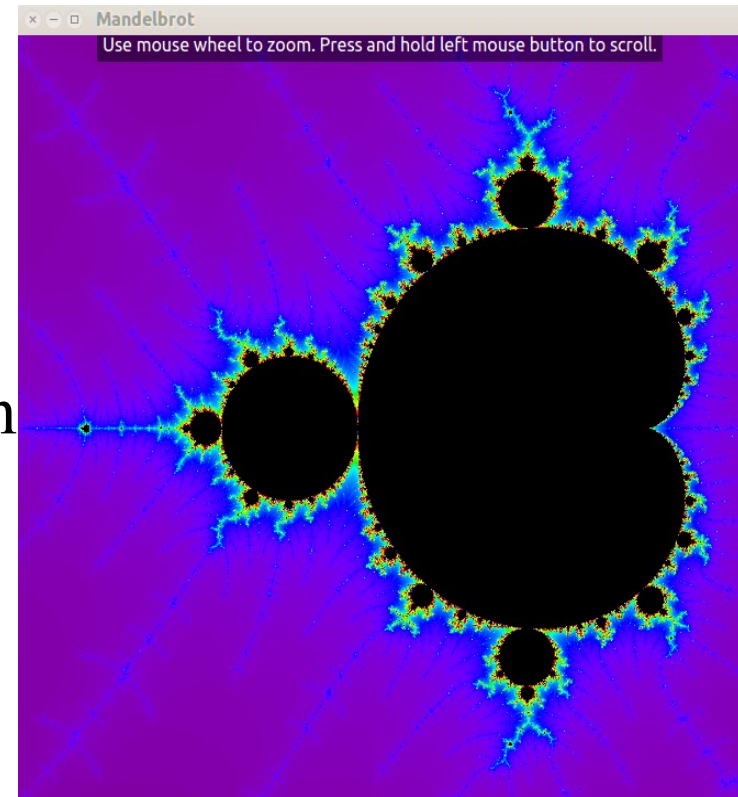


- The first stage reads a number of images from disk one by one, converts the images in B&W and produces a stream of images for the second stage
- The second stage applies the Sobel filter to each input image and then writes the output image into a separate disk directory

[Let's see the code!](#)

Examples: Mandelbrot set

- Very simple data-parallel computation
 - Each pixel can be computed independently
 - Simple ParallelFor implementation
- Black-pixel requires much more computation
- A naïve partitioning of the images quickly leads to load unbalanced computation and poor performance
 - Let's consider the minimum computation unit a single image line (image size 2048x2048, max 10^3 iterations per point)
 - **ParallelFor Static partitioning** of lines (48 workers) **MaxSpeedup 14**
 - **ParallelFor Dynamic partitioning** of lines (48 workers) **MaxSpeedup 37**

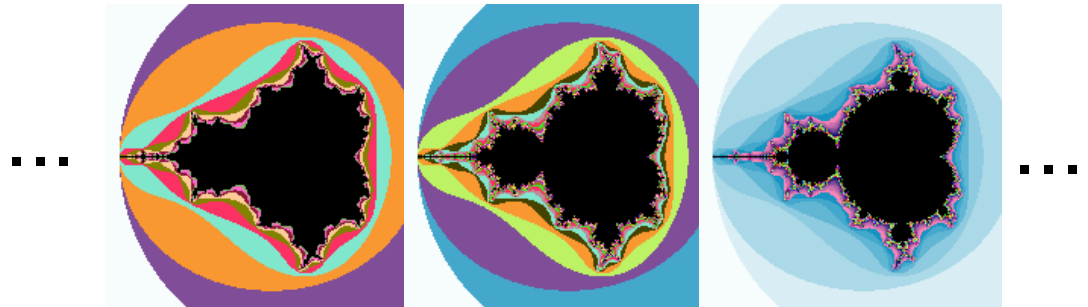


Data-partitioning may have a big impact on the performance

Examples: Mandelbrot set

(2)

- Suppose now we want to compute a number of Mandelbrot images (for example varying the computing threshold per point)



- We have basically two options:

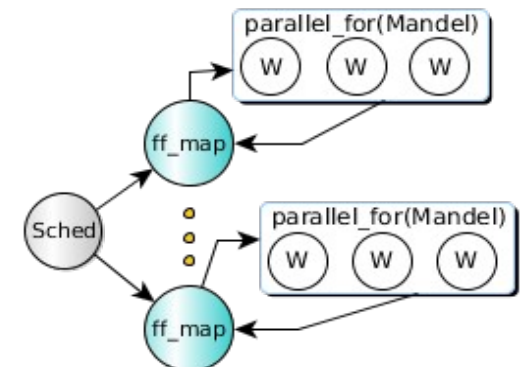
1. One single parallel-for inside a sequential for iterating over all different threshold points
2. A task-farm with map workers implementing two different scheduling strategies

```
for_each threshold values  
parallel_for ( Mandel(threshold));
```

- Which one is better having limited resources ?

– Depends on many factors, *too difficult to say in advance*

Moving quickly between the two solutions is the key point



Proposed exercises using ParallelFor & ParallelForReduce

- Finding the minimum and the index of the minimum value in an array of integer values of size N.

Given the following array:

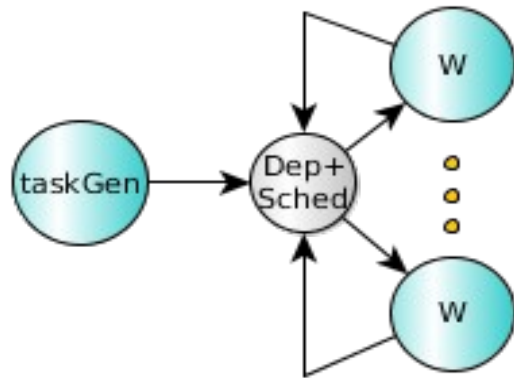
0	1	2	3	4	5	6	7	8	9
31	52	11	13	3	12	23	64	2	12

we want as output the pair $\langle 8, 2 \rangle$

- The sequential implementation is given in **arrayminindex.cpp**
- The objective is to implement the computation in parallel using the ParallelFor* patterns.

High-level patterns: ff_mdf

data-dependency pattern



```
void taskGen(ff_mdf*const mdf) {  
    ....  
    const param_info _1= {&A, INPUT};  
    const param_info _2 = {&B, INPUT};  
    const param_info _3 = {&C, OUTPUT};  
    std::vector<param_info> Param = {_1,_2,_3};  
  
    mdf->AddTask(Param, GEMM, A,B,C);  
    ....  
}
```

```
ff_mdf mdf(taskGen, ...);  
mdf.run_and_wait_end();
```

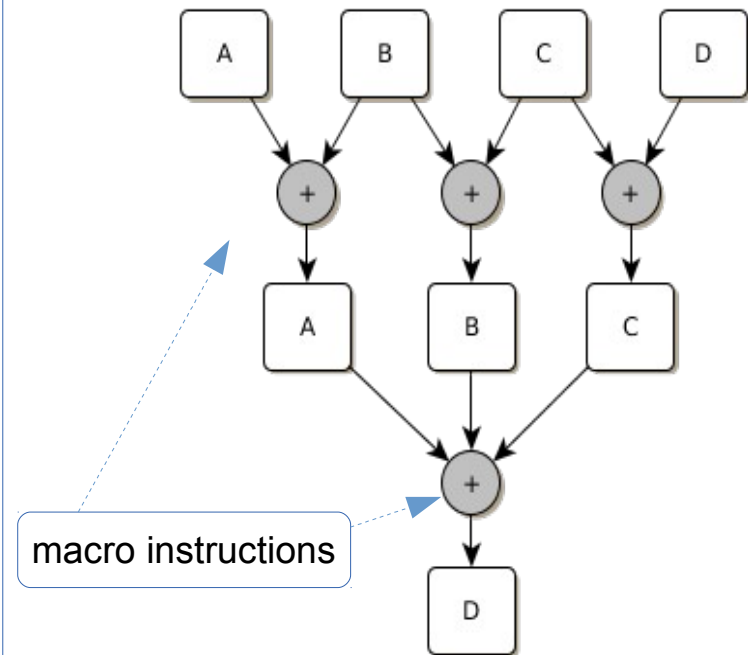
- The ff_mdf pattern targets *macro-data-flow* computations
- Is a general approach to parallelism based only on data dependencies.
- The computation is expressed by the data-flow graph, i.e. DAG whose nodes are macro-instructions and arcs are pure data-dependencies
 - A macro-instruction can be a set of simple instructions or a complex kernel function.
- By using the ff_mdf pattern, the user has to specify data-dependencies, i.e. declaring which are INPUT and OUTPUT data
- The **AddTask** method of the ff_mdf class is used to generate tasks
- The run-time, *automatically*, takes care of dependencies and then schedules ready tasks to Workers which executes ready (*fireable*) instructions in parallel

A simple test using ff_mdf

```
// macro operations
void sum2(long *X, long *Y, long size);
void sum3(long *X, long *Y, long *Z, long size);
....
// task generator function
void taskGen(Parameters<ff_mdf> *P) {
    ... auto mdf = P->mdf;

    { // A = A + B
        const param_info _1 = {&A, INPUT};
        const param_info _2 = {&B, INPUT};
        const param_info _3 = {&A, OUTPUT};
        std::vector<param_info> Param = {_1, _2, _3};
        mdf->AddTask(Param, sum2, A, B, SIZE);
    } { // B = B + C
        const param_info _1 = {&B, INPUT};
        const param_info _2 = {&C, INPUT};
        const param_info _3 = {&B, OUTPUT};
        std::vector<param_info> Param = {_1, _2, _3};
        mdf->AddTask(Param, sum2, B, C, SIZE);
    }
    .....
```

A,B,C,D are arrays of size N



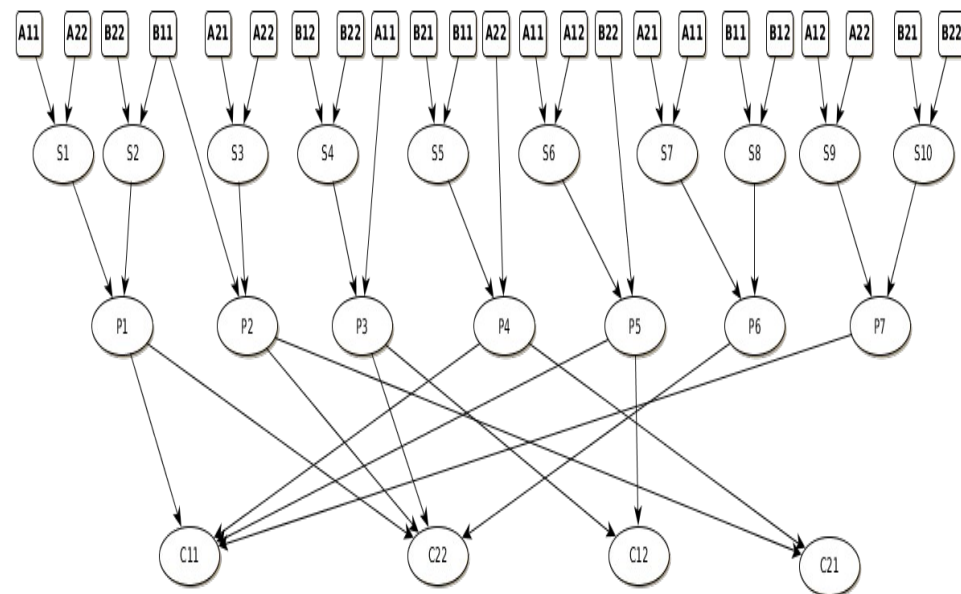
```
Parameters<ff_mdf> P; // structure containing all parameters needed to taskGen function
ff_mdf dag(taskGen, &P); // creates the mdf object
P.A=A,P.B=B,P.C=C....P.mdf=&dag... // preparing all parameters
dag.run_and_wait_end(); // run and wait termination
```

Strassen algorithm example using ff_mdf

- Matrix multiplication using Strassen's algorithm:
 - We want to compute $AxB = C$, A is $M \times P$, B is $P \times N$ and C is $M \times N$
 - Partitioning the matrices in 4 equal-size blocks we have:

$$\begin{array}{|c|c|} \hline A11 & A12 \\ \hline A21 & A22 \\ \hline \end{array} \times \begin{array}{|c|c|} \hline B11 & B12 \\ \hline B21 & B22 \\ \hline \end{array} = \begin{array}{|c|c|} \hline C11 & C12 \\ \hline C21 & C22 \\ \hline \end{array}$$

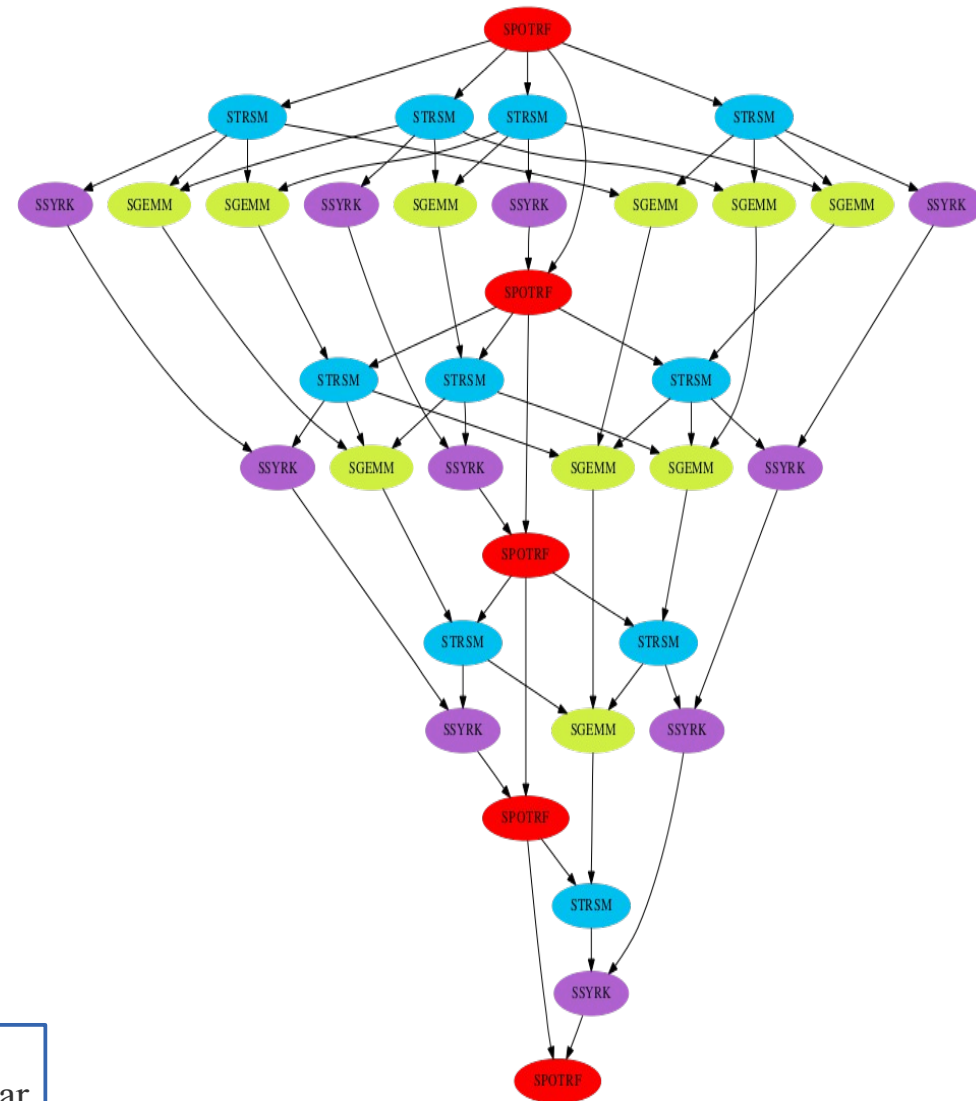
$$\begin{aligned} S1 &= A11 + A22 & S2 &= B11 + B22 & P1 &= S1 * S2 \\ S3 &= A21 + A22 & P2 &= S3 * B11 \\ S4 &= B12 - B22 & P3 &= A11 * S4 \\ S5 &= B21 - B11 & P4 &= A22 * S5 \\ S6 &= A11 + A12 & P5 &= S6 * B22 \\ S7 &= A21 - A11 & S8 &= B11 + B12 & P6 &= S7 * S8 \\ S9 &= A12 - A22 & S10 &= B21 + B22 & P7 &= S9 * S10 \\ C11 &= P1 + P4 - P5 + P7 \\ C12 &= P3 + P5 \\ C21 &= P2 + P4 \\ C22 &= P1 - P2 + P3 + P6 \end{aligned}$$



- Let's take a look at the code `examples/strassen/strassen_mdf.cpp` in the fastflow folder.

Block Cholesky and LU factorizations

- Much more complex data-flow graphs
- The DAG is huge for big matrices
 - cannot be entirely stored in memory
 - Have to be generated dynamically
- The *ff_mdf* pattern allows to generate the graph dynamically and overlap the graph generation phase with the computation



DAG represents, 5 tiles, left-looking ⁴³ version of Cholesky algorithm

D. Buono, M. Danelutto, T. De Matteis, G. Mencagli and M. Torquati “A light-weight run-time support for fast dense linear algebra on multi-core” in PDCN 2014 conference, 2014



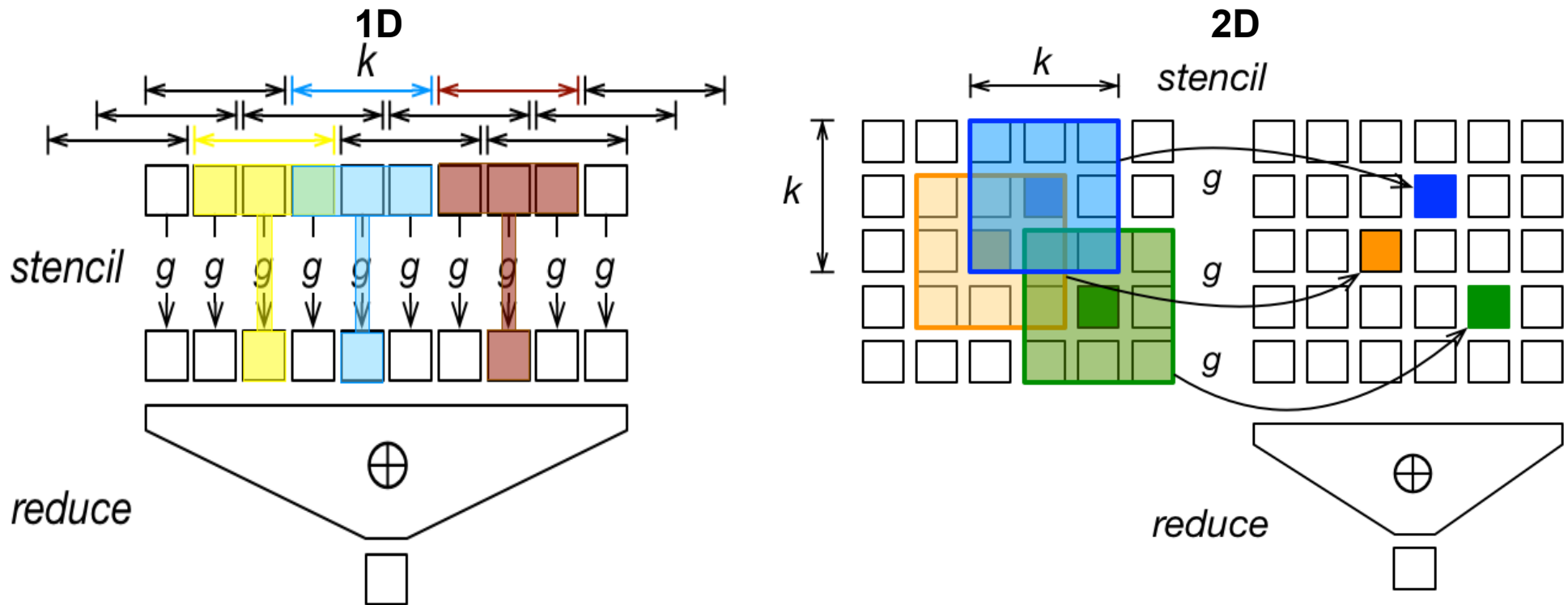
Targeting GPUs, basic concepts

(not yet covered in the FastFlow tutorial)

The Loop-Of-Stencil-Reduce pattern

- We introduced the **Loop-Of-Stencil-Reduce** pattern
 - Specifically targeting iterative Map+Reduce algorithm
 - Both 1D and 2D computation on multiple GPUs
 - On top of this pattern we implemented: map, reduce and map+reduce
 - OpenCL and CUDA version

Stencil computation



$$\text{stencil } g \ k \ a' = [g(S_0), g(S_1), \dots, g(S_{n-1})]$$

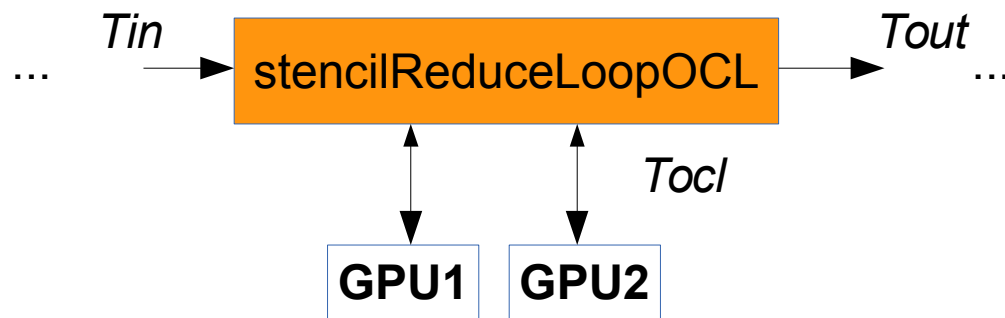
- g : list of $T \rightarrow$ list of T is the *elemental function*
- $S_i = [a'_{i-k}, \dots, a'_i, \dots, a'_{i+k}]$ is the *i-th neighborhood*
- we consider the case where k is a constant value
- a' is the extension of a with suitable extra items to deal with the border neighborhood computation

ff_stencilReduceLoop* run-time behaviour

```
// d2h input data, allocate double-buffering
(in_d_ptr, out_d_ptr, env_d_ptr) = allocate-write(input, env)
while cond
    // swap buffers, swap halo among different devices
    swap_buffers_and_halo (in_d_ptr, out_d_ptr)
    out_d_ptr = stencil <f_kernel> (in_d_ptr, env_d_ptr) // DEV
    part_d_ptr = reduce <op_kernel> (out_d_ptr) // DEV
    partial_data = read(part_d_ptr)
    cond = reduce op partial_data // host final reduction
// h2d-copy output, release buffers
output = release-read (in_d_ptr, out_d_ptr, env_d_ptr)
```

ff_stencilReduceLoop*

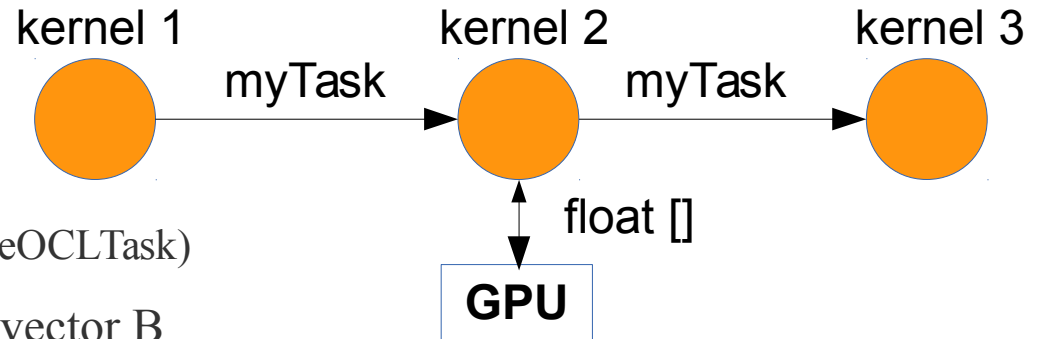
- Using OpenCL (for CUDA we have the same concepts)
- Implemented as a “special” kind of FastFlow node



- Each GPU is managed by a different internal “accelerator”
- `Tocl` is a template type encapsulating both input and output GPU types
 - It is used as an interface type via the `setTask/releaseTask` methods.
- Using the `ff_stencilReduceLoopOCL` we implemented:
 - `ff_mapOCL`, `ff_reduceOCL`, `ff_mapReduceOCL`
- Can be used as a pipeline stage or as a worker in a task-farm

Targeting GPUs in a pipeline

- Pipeline computation example:
 - Tin/Tout is myTask
 - Tocl is oclTask (must subclass from baseOCLTask)
- We want to compute a Map only on the vector B
 - B is non contiguous in memory



// this is the stream type (Tin/Tout)

```
struct myTask {  
  
    std::vector<float> A;  
    std::vector<std::vector<float> > B;  
    ....  
    std::string command;  
    ....  
};
```

```
// OpenCL interface type for the kernel2 node (Tocl)  
struct oclTask: baseOCLTask<myTask, float, float> {  
    // called by the run-time for each input task  
    void setTask(const Task *task) {  
        const size_t size = computeSize(B);  
        buffer = new float[size];  
        copyBuffer(B, buffer, size);  
        //set input and output host pointers  
        setInPtr(buffer, size); setOutPtr(buffer);  
    }  
    // called by the run-time at the end of task execution  
    void releaseTask(Task *task) {  
        copyBack(task->B, buffer); // copy results back to B  
        delete [] buffer;  
    }  
    float *buffer;  
};
```



Examples: Sobel filter

- Let's take a quick look at the `ff_pipe+mapOCL.cpp` example in the `fastflow/tests/ocl` folder.



Targeting distributed systems, basic concepts

(not yet covered in the FastFlow tutorial)

Targeting distributed systems (concepts)



- ***Building blocks:***

- **ff_dnode** extend the **ff_node** class

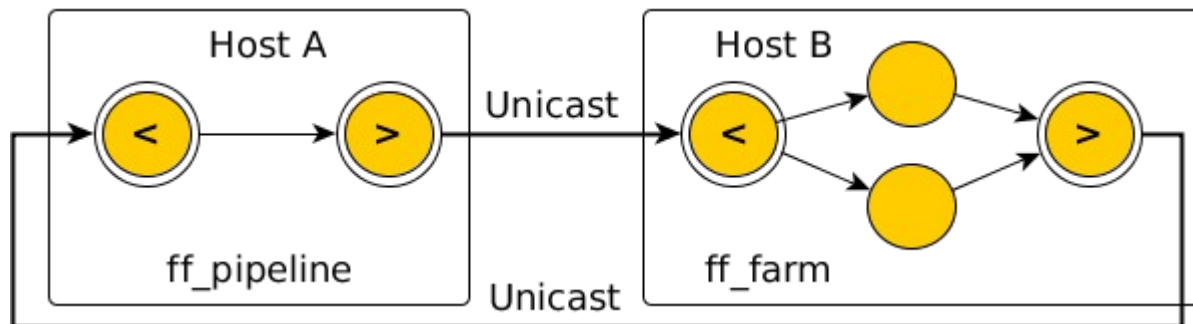
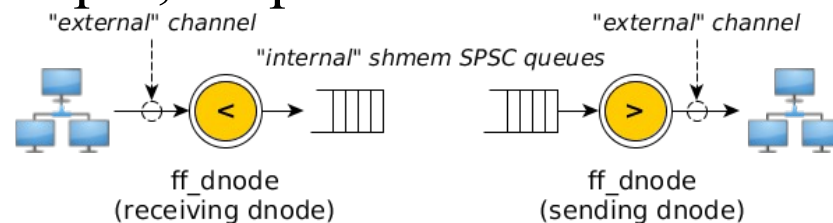
- **Communication patterns:**

- **Unicast: one-to-one** unidirectional channel
- **Broadcast: one-to-many**, the same data is sent to all connected peers
- **Scatter: one-to-many**, data is split in disjoint partitions, each one sent to a distinct connected peer
- **On-demand: one-to-one**, the data is sent to one of the connected peer using a request-reply protocol
- **Gather-All: many-to-one**, receives a data partition from all connected peers, and eventually the message is recomposed
- **Collect-from-Any: one-to-one**, receives data from one of the connected peers
- **TODO: Multicast,**

-

ff_dnode

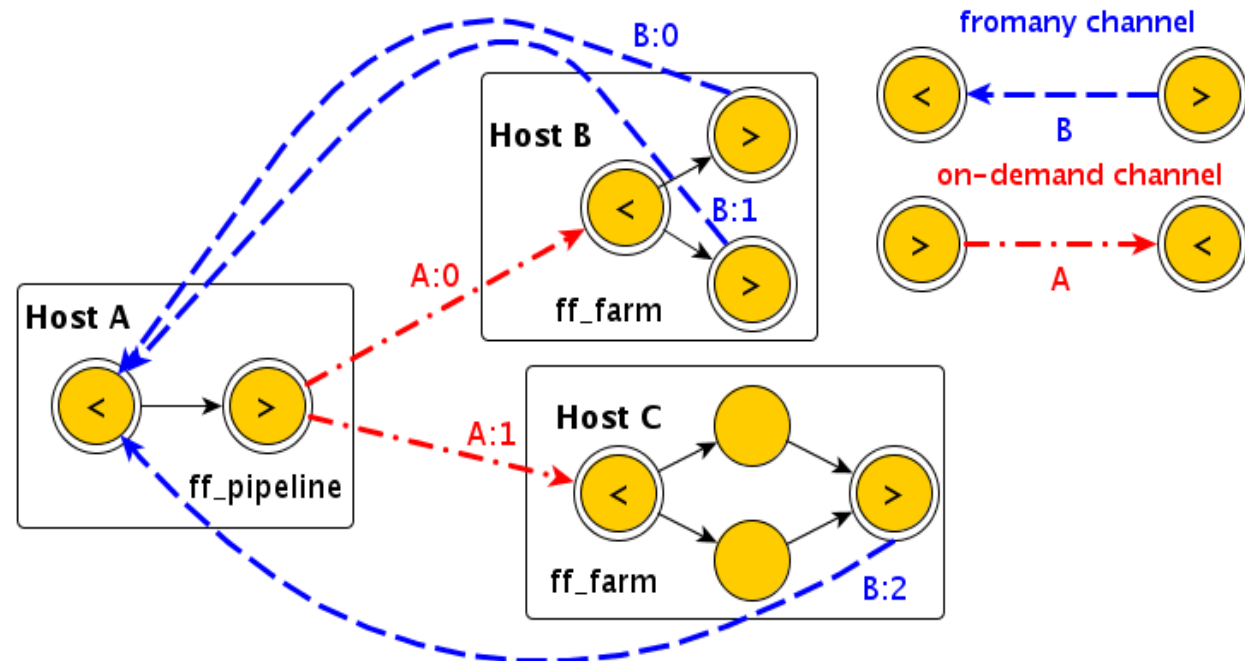
- A `ff_dnode` is an `ff_node` with an extra channel used to communicate with another dnode by using a communication pattern (external channel)
- The external channel may be specialized for input, output or both
- The idea is that *only edge nodes* are dnodes



- Communication patterns:
 - TCP/IP networks by using the ZeroMQ library
 - Infiniband networks by using Linux verbs (experimental)

How it works

- The user has to:
 - Define and connect the application parts using proper communication patterns
 - Prepare the data to send
 - Provide the run-time with message buffers where data has to be received



How to define a dnode

```
emacs@pirotecni
File Edit Options Buffers Tools C++ Help
// just an edge-node of my application
class myNode: ff_dnode<zmqBcast> {
public:
    myNode(const std::string& name,          // unique identifier for the channel
           const std::string& address,      // peer address [host:port]
           const int npeers,                // n. of peers
           zmqBcast::TransportImp* const t); // transport object

    int svc_init() {
        // I am the producer, i.e. the one that sends data out.
        // The consumer has to call init with the last param set to false.
        return ff_dnode<zmqBcast>::init(name,address,npeers,t, true);
    }

    void* svc(void* task) {
        // As soon as a task is returned or the ff_send_out method is called,
        // the message is sent out in broadcast (because of zmqBcast pattern)
        // to all connected peers.
    }

    void svc_end() { ....}
};
```



Distributed image filtering (img_d.cpp)

- Let's take a quick look at the `img_d.cpp` example which is a possible distributed version of the `img.cpp` app.
 - This version uses only 2 machines
 - It is just a proof-of-concept implementation

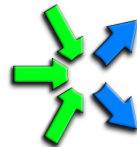
GridKa School 2015

Big Data Virtualization Modern Programming

That's (almost) all !
Thanks for participating!

For any questions or comments please send an e-mail to torquati@di.unipi.it

FastFlow project site:



<http://mc-fastflow.sourceforge.net>

<http://calvados.di.unipi.it/fastflow>

EU-FP7 projects using FastFlow:

